

## EAST Search History

Ref #	Hits	Search Query	DBs	Default Operator	Plurals	Time Stamp
L1	4	"672811".ap.	US-PGPUB; USPAT	OR	OFF	2007/10/11 21:12
L2	1476	(380/28).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 21:49
L3	1157	2 and @ay <="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 21:49
L4	11	3 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 22:09
L5	0	(derivative) same ("not" determin\$4 factor)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/10/11 21:54
L6	0	false derivate	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	OFF	2007/10/11 21:54
L7	0	false derivative value	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/10/11 21:54

## EAST Search History

L8	0	(derivative value) same ("not" determin\$4 factor)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/10/11 21:55
L9	0	(derivative value) same ("not" with factor)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/10/11 21:55
L10	121	(380/263).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 21:55
L11	87	10 and @ay<="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 21:56
L12	5	11 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 21:56
L13	1339	(380/30).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 21:57
L14	1179	13 and @ay<="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 21:57

## EAST Search History

L15	7	14 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 21:57
L16	806	(380/46).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 21:57
L17	720	16 and @ay <="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 21:57
L18	8	17 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 21:58
L19	350	(380/42).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 21:59
L20	303	19 and @ay<="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 21:59
L21	1	20 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 22:02

## EAST Search History

L22	261	(380/45).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:02
L23	249	22 and @ay <="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:02
L24	1	23 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 22:03
L25	395	(713/180).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:03
L26	352	25 and @ay<="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:03
L27	1	26 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 22:04
L28	2015	(713/168).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:04

## EAST Search History

L29	1476	28 and @ay<="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:04
L30	6	29 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 22:05
L31	281	(708/491).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:05
L32	257	31 and @ay<="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:06
L33	6	32 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 22:07
L34	308	(708/492).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:08
L35	265	34 and @ay<="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:08

## EAST Search History

L36	5	35 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 22:08
L37	5410	(709/217).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:09
L38	4306	37 and @ay<="2003"	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:09
L39	12	38 and (factor) and (equation) and (derivative)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/10/11 22:09
L40	0	((encrypt\$4) and (decrypt\$4) and (original string) and (defin\$4 set of factors) and (encrypt\$4 equation) and (map\$4) and (a set of derivative equations) and (generat\$4 derivative values) and (stor\$4 encrypted string and derivative values) and (provid\$4 false derivatives cannot be used to determine a given factor) and (stor\$4 false derivative values with the generated derivative values) and (us\$4 a set of factor decryption equatins to map) and (decrypt\$4 the encrypted string) and (decryption equation) and (factor mapped through the set of factor decryption equation) and (presence) and (false derivative values with the generated derivative values) and (prevent\$4) and (use with the factor decryption equation to derive the factors in the set of facotrs)).clm.	US-PGPUB; USPAT	ADJ	ON	2007/10/11 22:47

## EAST Search History

L41	6	Cheung-tom-thuan.in.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/10/11 22:48
S1	2	"672811".ap.	USPAT	OR	OFF	2007/03/02 10:44
S2	3	"672811".ap.	US-PGPUB; USPAT	OR	OFF	2006/12/01 16:29
S3	194	( "20050069127" "4893338" "5857025" "6049608" "6078665" "20020154769" "20040086117" "5666419" "6259790" "6317769" "20020055962" "20030081769" "20030215089" "20050195973" "6058189" "4306111" "5675653" "4969190" "5727064" "5982900" "6128386" "6330674" "6449473" "7133522" "20030007639" "20030059040" "20030118185" "20030210781" "4423287" "4578530" "5313521" "5214698" "5953419" "5265164" "5517567" "5796830" "5987133" "6003136" "6009175" "6055639" "6175920" "6230002" "5224166" "5365589" "5412729" "5481613" "5539827" "5694470" "5787175" "5799086" "5815573" "5841865" "5850451" "5857022" "5864683" "5872849" "5917910" "5920630" "5937066" "5956408" "6009177" "6052469" "6061454" "6061454" "6072876" "6122742" "6185316" "6202150" "6249866" "4295039" "4302810" "4438824" "4500750" "4817140" "4825050" "4881263" "4924513" "4972472" "5001754" "5010573" "5200999" "5208853" "5222140" "5228084" "5295188" "5345506" "5365588" "5384850" "5420866" "5425103" "5434920" "5448638" "5455861" "5495533" "5541994" "5594798" "5604801" "5615268" "5619576" ).pn.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2006/12/01 16:54

## EAST Search History

S4	24	("5677952" "5751811" "5781632"  "5805712" "5987133" "6002769"  "6226618" "6230272" "6324287"  "6345288" "6351813" "6389541" ).PN.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2006/12/01 16:56
S5	1	"5862325".pn.	USPAT	OR	OFF	2006/12/01 16:55
S6	1	"5915024".pn.	USPAT	OR	OFF	2006/12/01 16:55
S7	1	"6226383".pn.	USPAT	OR	OFF	2006/12/01 16:55
S8	40	(derivative equation) and (factor) and (encrypt\$4) and (decrypt\$4)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2006/12/04 20:32
S9	1287	(380/28).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2006/12/04 20:33
S10	2	S9 and S8	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2006/12/04 20:33
S11	44	("4074066"   "4238854"   "4319079"   "4593353"   "4734796"   "4888798"   "4907274"   "5003597"   "5212729"   "5239581"   "5454039").PN. OR ("5677952"). URPN.	US-PGPUB; USPAT; USOCR	OR	OFF	2006/12/06 13:36
S12	2583	(password) and (deriv\$5) and (factor) and (encryp\$4) and (decrypt\$4)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2006/12/07 13:57
S13	934	(password) and (deriv\$5) and (factor) and (encryp\$4) and (decrypt\$4)	USPAT	OR	ON	2006/12/07 13:58



## EAST Search History

S14	5	(password) same (deriv\$5) same (factor) same (encryp\$4) same (decrypt\$4)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2006/12/07 16:00
S15	1	"6845453".pn.	USPAT	OR	OFF	2006/12/07 14:01
S16	13	("5280527"   "5499297"   "5805719"   "5907597"   "6035398"   "6035406"   "6307955"   "6317834"   "6332193"   "6385318"   "6553494").PN. OR ("6845453"). URPN.	US-PGPUB; USPAT; USOCR	OR	OFF	2006/12/07 14:05
S17	1	(password) same (derivative) same (factor) same (encryp\$4) same (decrypt\$4)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2006/12/07 15:26
S18	900	(password) and (derivative) and (factor) and (encryp\$4) and (decrypt\$4)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2006/12/07 15:26
S19	341	(password) and (derivative) and (factor) and (encryp\$4) and (decrypt\$4)	USPAT	OR	ON	2006/12/07 15:27
S20	1274	(380/28).CCLS.	US-PGPUB; USPAT; USOCR; FPRS; JPO; DERWENT; IBM_TDB	OR	OFF	2006/12/07 15:27
S21	6	S18 and S20	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2006/12/07 15:31

## EAST Search History

S22	2390	factor same encryp\$4	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2006/12/07 15:32
S23	956	factor same encryp\$4	USPAT	OR	ON	2006/12/07 15:32
S24	360	factor with encryp\$4	USPAT	OR	ON	2006/12/07 15:32
S25	1	"6075865".pn.	USPAT	OR	OFF	2006/12/07 16:00
S26	3	("5375169"   "5787173"   "5835597").PN. OR ("6075865"). URPN.	US-PGPUB; USPAT; USOCR	OR	OFF	2006/12/07 16:07
S27	1	"5048086".pn.	USPAT	OR	OFF	2006/12/08 09:38
S28	1	"20020176377".pn.	US-PGPUB; USPAT	OR	OFF	2007/03/02 10:45
S29	3	"672811".ap.	US-PGPUB; USPAT	OR	OFF	2007/05/15 09:58
S30	3	"672811".ap.	US-PGPUB; USPAT	OR	OFF	2007/05/23 11:23
S31	1	"5048086".pn.	US-PGPUB; USPAT	OR	OFF	2007/05/23 11:23
S32	62	("5007087").PN. OR ("5048086"). URPN.	US-PGPUB; USPAT; USOCR	OR	OFF	2007/05/23 11:34
S33	0	generat\$4 with (random number) with ("not") with chaos	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:36
S34	0	("not") with chaos	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:35
S35	1	("not") same chaos	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:35

## EAST Search History

S36	0	generat\$4 same (random number) same ("not") same chaos	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:36
S37	0	(random number) same ("not") same chaos	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:37
S38	10	chaos equation	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:44
S39	0	(encrypt\$4) with ("not") with (base\$2) with (chao\$4)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:41
S40	0	(encrypt\$4) same ("not") same (base\$2) same (chao\$4)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:41
S41	4282	("not" or "no" or "unlike") same (chaos)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:45
S42	2265	("not" or "no" or "unlike") with (chaos)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:45

## EAST Search History

S43	2256	("not" or "no") with (chaos)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:48
S44	168	(without) with (chaos)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:45
S45	154	S42 and (encrypt\$4)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/23 11:47
S46	168	("not" or "without") with (chaos)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 11:49
S47	17	S46 and (encrypt\$4)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	ADJ	ON	2007/05/23 12:02
S48	0	"672811".pn.	US-PGPUB; USPAT	OR	OFF	2007/05/23 15:16
S49	3	"672811".ap.	US-PGPUB; USPAT	OR	OFF	2007/05/23 15:16
S50	1	"5677952".pn.	USPAT	OR	OFF	2007/05/24 10:03
S51	1295	(380/28).CCLS.	US-PGPUB; USPAT	OR	OFF	2007/05/24 12:16
S52	0	("672811.ap.").PN.	US-PGPUB; USPAT; USOCR	OR	OFF	2007/10/06 20:43
S53	4	"672811".ap.	US-PGPUB; USPAT	OR	OFF	2007/10/06 20:43

# ETC Comprehensive search

File 696:DIALOG Telecom. Newsletters 1995-2007/Oct 10  
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S1	1367162	FALSE OR INVALID OR NULL OR PSUEDO OR PSEUDO OR RANDOM OR - RAND OR CHAFF OR SEMIRANDOM OR PSUEDORAND? OR PSEUDORAND? OR - DUMMY OR NONCE? ? OR DECOY?
S2	196815	SHILL?? OR SPOOF?? OR PHONEY? ? OR PHONY? ? OR PHONIE? ? - OR FAKE? ? OR SHAM OR SHAMS
S3	31249	S1:S2(1W)(VALUE OR VALUES OR NUMBER? ? OR NUMERAL? ? OR NU- MERIC?? OR ALPHANUMERIC? OR QUANTITY? ? OR QUANTITIES OR CHAR- ACTER? ? OR BIT OR BITS OR DATA OR DATUM? ?)
S4	82185	S1:S2(1W)(DIGIT? ? OR SIGNAL? ? OR PULSE OR PULSES OR INTE- GER? ? OR STRING OR STRINGS OR SUBSTRING? ? OR SEQUENCE OR SE- QUENCES OR SUBSEQUENCE? ? OR NOISE OR VARIABLE? ?)
S5	29552231	USED OR USE OR USING OR USAGE? ? OR EMPLOY? OR UTILIS??? OR UTILISATION? ? OR UTILIZ??? OR UTILIZATION? ?

S6 489797 NONEMPLOY? OR UNEMPLOY? OR NONUTILIS? OR NONUTILIZ? OR UNU-  
TILIS? OR UNUTILIZ?  
S7 142283 NONUSE? OR NONUSING OR NONUSAGE? OR UNUSE? OR UNUSING OR U-  
NUSAGE? OR DISUSE? ? OR DISUSING OR DISUSAGE?  
S8 1035688 ('NOT' OR WITHOUT OR NEVER OR CANNOT OR T) (1W)S5  
S9 18436 (NON OR UN OR DIS)()S5  
S10 1442 S3:S4(5N)(OMIT? OR OMISSION? ? OR EXCLUD? OR EXCLUSION?)  
S11 163 S3:S4(5N)S8:S9  
S12 261023 DECRYPT? OR UNCRYPT? OR UNENCRYPT? OR DECOD??? OR UNENCOD?-  
?? OR UNCOD??? OR UNENCRYPT? OR UNENCIPHER? OR UNENCYPHER? OR  
DECIPHER? OR DECPHER?  
S13 99025 UNSCRAMBL? OR DESCRAMBL? OR UNCIPHER? OR UNCYPHER? OR DECO-  
MPRESS? OR DEPACK? OR UNPACK? OR UNCOMPRESS? OR DECOMPACT? OR  
UNCOMPACT?  
S14 214 S3:S4(5N)S6:S7  
S15 3 (S10:S11 OR S14)(50N)S12:S13  
S16 1 RD (unique items)  
S17 211 AU=(CHEUNG T? OR CHEUNG, T?)  
S18 0 S17 AND (S10:S11 OR S14)  
S19 0 S17 AND S3:S4  
S20 7 (S10:S11 OR S14)(100N)S12:S13  
S21 4 S20 NOT S15  
S22 2 RD (unique items)

22/3,K/1 (Item 1 from file: 484)  
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Prying open the Clipper lock

Anonymous

Science News (GSCN), v145 n24, p383, p.1

Jun 11, 1994

ISSN: 0036-8423 JOURNAL CODE: GSCN

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WORD COUNT: 331

LENGTH: Short (1-9 col inches)

TEXT:

... it by taking advantage of the fact that to obtain the right master key to decipher a message, the government has to find out the encryption chip's serial number. This...

...legitimate. It then substitutes a bogus LEAF, which has the same checksum but includes a fake serial number that can't be used by the government. Officials of the National Security Agency concede the existence of the flaw...

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02368408 SUPPLIER NUMBER: 02743601 (USE FORMAT 7 OR 9 FOR FULL TEXT)

Basic cryptography: SBOEPN DJQIFST.

Wheeler, Daniel D.; Perri, Elisheva

Creative Computing, v9, p178(4)

May, 1983

ISSN: 0097-8140

LANGUAGE: ENGLISH

RECORD TYPE: FULLTEXT

WORD COUNT: 2134

LINE COUNT: 00156

... 130 is included only to make a complete statement; the value stored in X is never used .

Setting the random number seed on the IBM PC is a trivial process, as the RANDOMIZE function allows automation...also be noted that messages for enciphering must be input in upper case, for proper decoding. The rest of the program follows other Microsoft versions closely. There is no instruction to...

22/9/2 (Item 1 from file: 47)  
DIALOG(R)File 47:Gale Group Magazine DB(TM)  
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02368408 SUPPLIER NUMBER: 02743601 (THIS IS THE FULL TEXT)  
Basic cryptography: SBOEPN DJQIFST.  
Wheeler, Daniel D.; Perri, Elisheva.  
Creative Computing, v9, p178(4)  
May, 1983  
ISSN: 0097-8140 LANGUAGE: ENGLISH RECORD TYPE: FULLTEXT  
WORD COUNT: 2134 LINE COUNT: 00156

TEXT:

Basic Cryptography: SBOEPN DJQIFST

The personal computer is a powerful tool for cryptography. With a bit of simple programming you can encipher secret messages to your friends so securely that it would take the efforts of a mathematical cryptographer to unravel the system. But with the proper keyword a friend can use his computer to decipher and read your message.

There have been several programs for simple cryptography published in the major computer magazines. Some of these programs, unfortunately, have used very weak systems for enciphering the messages. Anyone with a little knowledge of cryptanalytic techniques could break the system and read the messages without knowing the keyword.

This article explains the principles of computer cryptography and demonstrates the use of the Basic random number function for enciphering messages. Versions of the program are included for the Apple II, the TRS-80 Models I and III, IBM PC, and the Atari 400 and 800.

Cryptography by Addition

Letters are represented in computers as numbers. This makes it easy to use the arithmetic operations of Basic to transform the letters. The simplest system is just to add a constant number to the character code for each letter. If the result is too large, subtract the number of characters being used so that the result is again a valid character code. Suppose, for instance, that you want to use three as the constant and that your messages consist of only capital letters. Then each letter in the message will come out as the letter three letters further on in the alphabet. The letter A (ASCII code 65) will appear as D (ASCII code 68), B(66) becomes E(69), and so forth. At the end of the alphabet, X (88) becomes the ASCII code 91. But 90 is Z and 91 is beyond the end of the alphabet. So 26 (the number of characters we are using) is subtracted from 91 to produce 65. Thus X wraps around to the beginning of the alphabet and becomes A.

The Basic statements necessary to do this are quite simple. If the letter to be transformed is stored in the string A\$, you can do it with:

```
100 = ASC(A$ )
110 X = X + 3
120 IF > 90 THEN = - 26
130 A$ = CHR$ (X)
```

The ASC(A\$ ) function converts the character to a numeric variable so that the arithmetic can be done in the next two statements. The CHR\$ (X) function converts the numeric result back into a string.

To decipher the message, change lines 110 and 120 to:

```
110 X = X - 3
120 IF < 65 THEN = 26
```

This system is called a Caesar cipher because Julius Caesar is said to have used it. It may have fooled the Gauls, but now any bright elementary school student (maybe with a hint) can break the system. Part of

the title of this article is in the Caesar cipher, but not with an offset of three.

#### Better Systems

The weakness of the system comes from the use of the constant. There are only 25 possible constants to try and once you figure it out it is easy to decipher the whole message. You can improve the system by changing the offset for each character. You might try adding one to the offset each time you encipher a character and then subtracting 26 from the constant whenever it gets too large. This will produce a cipher that is much more difficult to break.

There are many possible schemes for changing the offset. Any scheme will work to make the cipher more difficult to break. But if the scheme is simple (like adding one) and repeats at fairly short, regular intervals, then it is not very difficult to figure out the pattern and break the cipher. A smart high school student could do it.

What you need is an irregular pattern that doesn't repeat within the length of the messages you are interested in sending. The random number function in Basic provides a very irregular series of numbers. They do repeat eventually, but the cycle is much longer than any messages you will ever send on your computer.

If the random number function were truly random, it would not be useful for our purpose. Once you have enciphered a message, your recipient must be able to generate the same sequence of numbers to decipher the message. Fortunately, most versions of Basic provide some way to "seed" the random function so that it generates the same sequence of numbers.

In Applesoft, for instance, calling the random number function with a negative argument, such as RND (-99), seeds the generator to start at a definite place in the sequence. If you agree beforehand on a number to use as the seed, your friend will be able to decipher your message by generating the same sequence of numbers to use as offsets.

#### Demonstration Programs

Listings 1, 2, 3, and 4 show programs to demonstrate these techniques for four popular microcomputers. Each of the programs enciphers or deciphers a one-line secret message. Instead of enciphering just the letters of the message, these programs encipher everything: letters, numbers, punctuation marks and even spaces. (The ASCII code for the space is 32. It is just as much a character as any of the others. You must be especially careful in typing the enciphered message to get all the spaces exactly right.)

The program lines in the 100's initialize the random number generator. For the Apple this is simply a matter of calling the random number generator with a negative argument. The variable X in line 130 is included only to make a complete statement; the value stored in X is never used.

Setting the random number seed on the IBM PC is a trivial process; as the RANDOMIZE function allows automation of the seed generation. By omitting an argument in the RANDOMIZE command in line 110, the PC will return with the default input statement, Random Number Seed (-32768 to 32767)? You may then input your cipher base value. The message is input with an INKEY command, so backspacing is impossible. It should also be noted that messages for enciphering must be input in upper case, for proper decoding. The rest of the program follows other Microsoft versions closely.

There is no instruction to initialize the RND() function in TRS-80 Level II Basic, but it can be done with POKES into memory. Lines 130-150 show how to do it. POKES can only be done with numbers smaller than 256. The instructions in lines 140 and 150 break the larger seed (stored in N) into two parts, each less than 256.

We couldn't figure out how to seed the RND() function in Atari Basic, so we'll show you how to write your own random function. The initial seed must be a decimal fraction between zero and one. In lines 110-140 the program gets a number and then divides by 100,000 to make it a fraction.

The lines in the 200's allow you to select whether the message will be enciphered (by adding the random numbers) or deciphered (by subtracting



the random numbers).

The next section of the program (300's) allows you to enter your message. For the Atari this is a straight-forward INPUT statement. Then the loop in lines 340-360 converts the characters to the numeric (ASCII code) values and stores them in the array IN(). But neither the Apple nor the TRS-80 allows commas within input strings. The comma is used to separate multiple items in the input. Since we wanted to include the comma as an allowable character we used the single character input commands. These are GET on the Apple and INKEY\$ on the TRS-80. The program loop starting at line 330 accepts single characters, converts them to numeric form, and stores the ASCII codes in the integer array IN%().

When the message is completely entered, the program goes to the section either to encipher (400's) or decipher (500's) the message. There are 59 possible characters from 'space' (ASCII 32) to Z (ASCII 90). To encipher the message we should add a random integer up to 59 to each of the character codes. This is easy on the TRS-80.

The RND() function with arguments larger than one returns integers in the range from one to the value of the argument. Thus RND(59) returns integers from 1 to 59. These are added to the character codes in line 420. Line 430 subtracts 59 if the result is out of the allowable range. Line 430 converts the numeric code to a character and prints it. The loop in lines 410-450 repeats this for each character in the message.

The Apple RND() function returns decimal fractions between zero and one. To convert to a random integer we multiply by 59 and use the INT() function to make the result an integer. This appears in line 420. The rest of the loop is exactly the same as for the TRS-80.

In the Atari version we don't use the built-in RND() function. We store the seed for our own random function in the variable N. To get each successive random number we multiply N by 997 and take the fractional part of the result to use as the random number and to store in N for generating the next number. Line 420 does this by calculating  $997*N$  and subtracting the integer part to leave the fractional part. Then N is used in line 430 as a random number in the range zero to one, just as in the Apple version.

The section to decipher the message (lines in the 500's) is exactly the same as the enciphering section except that the additions and subtractions are reversed. It will restore an enciphered message to its original form.

#### Extending the Demonstration Programs

These demonstration programs are not intended for practical use. They can, however, be extended to meet your cryptographic needs. You will certainly want to put in a loop so that your messages can be more than one line long. You will probably want the output written on disk or cassette so the person receiving your message won't have to type the random-appearing enciphered text. Output to a modem for telephone communication is another possibility. Your imagination is the only limit.

#### Breaking Random Ciphers

You might think that the ciphers based on random number generators would be impossible to break. After all, the enciphered message looks just like a random sequence of characters. There is no pattern to give clues to the content of the message. During World War II the Germans were confident that their machine cipher was secure. But first the Poles and then the British were able to break it. Churchill was reading Hitler's war dispatches--sometimes even before they got to Hitler.

The method requires that the cryptographer be able to guess a word in the message. For instance, if the message looks as though it was intended as a letter, it is likely to begin "Dear . . ." The cryptographer subtracts the ASCII codes for "Dear" from the message to recover part of the sequence of the random number generator. It is possible to figure out from a few numbers where the random number generator is in its sequence. Then it is a simple matter to generate the entire sequence and decipher the whole message. If the first attempt doesn't work, the cryptographer tries other probable words in all possible positions in the message.

There are techniques for enciphering messages that are resistant to the probable word method. If you have a serious security problem you should

get a commercially available, tested system. But for most personal computer users the ciphers based on the Basic random function provide a reasonable degree of security. Unless your lover's spouse is a mathematician, you'll be able to keep your letters secret with Basic random ciphers.

Table: Listing 1. Apple II version of the random cipher program.

Table: Listing 2. TRS-80 version of the random cipher program.

Table: Listing 3. Atari version of the random cipher program.

Table: Listing 4. IBM PC version of the random cipher program.

Table: Sample Run.

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DESCRIPTORS: cryptography--Equipment and supplies; Ciphers--Computer programs

FILE SEGMENT: MI File 47

File 348:EUROPEAN PATENTS 1978-2007/ 200739

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File 349:PCT FULLTEXT 1979-2007/UB=20070927UT=20070920

(c) 2007 WIPO/Thomson

Set	Items	Description
S1	477716	FALSE OR INVALID OR NULL OR PSUEDO OR PSEUDO OR RANDOM OR - RAND OR CHAFF OR SEMIRANDOM OR PSUEDORAND? OR PSEUDORAND? OR - DUMMY OR NONCE? ? OR DECOY?
S2	13399	SHILL?? OR SPOOF?? OR PHONEY? ? OR PHONY? ? OR PHONIE? ? - OR FAKE? ? OR SHAM OR SHAMS
S3	34577	S1:S2(1W)(VALUE OR VALUES OR NUMBER? ? OR NUMERAL? ? OR NU- MERIC?? OR ALPHANUMERIC? OR QUANTITY? ? OR QUANTITIES OR CHAR- ACTER? ? OR BIT OR BITS OR DATA OR DATUM? ?)
S4	31600	S1:S2(1W)(DIGIT? ? OR SIGNAL? ? OR PULSE OR PULSES OR INTE- GER? ? OR STRING OR STRINGS OR SUBSTRING? ? OR SEQUENCE OR SE- QUENCES OR SUBSEQUENCE? ? OR NOISE OR VARIABLE? ?)
S5	2256842	USED OR USE OR USING OR USAGE? ? OR EMPLOY? OR UTILIS??? OR UTILISATION? ? OR UTILIZ??? OR UTILIZATION? ?
S6	1304	NONEMPLOY? OR UNEMPLOY? OR NONUTILIS? OR NONUTILIZ? OR UNU- TILIS? OR UNUTILIZ?
S7	40445	NONUSE? OR NONUSING OR NONUSAGE? OR UNUSE? OR UNUSING OR U- NUSAGE? OR DISUSE? ? OR DISUSING OR DISUSAGE?
S8	397884	('NOT' OR WITHOUT OR NEVER OR CANNOT OR T) (1W)S5
S9	24360	(NON OR UN OR DIS)()S5
S10	376	S3:S4(5N)(OMIT? OR OMISSION? ? OR EXCLUD? OR EXCLUSION?)
S11	409	S3:S4(5N)S8:S9
S12	146747	DECRYPT? OR UNCRYPT? OR UNENCRYPT? OR DECOD??? OR UNENCOD?- ?? OR UNCOD??? OR UNENCRYPT? OR UNENCIPHER? OR UNENCYPHER? OR DECIPHER? OR DECYPHER?
S13	45013	UNSCRAMBL? OR DESCRAMBL? OR UNCIPHER? OR UNCYPHER? OR DECO- MPRESS? OR DEPACK? OR UNPACK? OR UNCOMPRESS? OR DECOMPACT? OR UNCOMPACT?
S14	119	S3:S4(5N)S6:S7
S15	77	(S10:S11 OR S14)(50N)S12:S13
S16	56	S15 AND PY=1963:2003
S17	31	S15 AND (AC=US OR AC=US/PR) AND AY=1978:2003
S18	61	S16:S17
S19	61	IDPAT (sorted in duplicate/non-duplicate order)
S20	61	IDPAT (primary/non-duplicate records only)

? t20/5,k/7,16,19,21,27,31,36,38-39,50-51,61

20/5,K/7 (Item 7 from file: 348)

DIALOG(R)File 348:EUROPEAN PATENTS

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01439227

Method, apparatus and program for quantitative competition and recording  
medium having recorded thereon the program

Verfahren, Vorrichtung und Programm für quantitative Wettkämpfe und  
Aufzeichnungsmedium, welches das Programm gespeichert hat

Systeme, methode et programme pour la concurrence quantitative et medium de  
stockage pour le programme associe

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PATENT (CC, No, Kind, Date): EP 1225530 A1 020724 (Basic)  
EP 1225530 B1 051026

APPLICATION (CC, No, Date): EP 2002001009 020117;

PRIORITY (CC, No, Date): JP 200110327 010118

DESIGNATED STATES: DE; FR; GB

EXTENDED DESIGNATED STATES: AL; LT; LV; MK; RO; SI

INTERNATIONAL PATENT CLASS (V7): G06F-017/60

CITED REFERENCES (EP B):

KIKUCHI H.; HOTTA S.; ABE K.; NAKANISHI S.: 'distributed auction servers  
resolving winner and winning bid without revealing privacy of bids'

PARALLEL AND DISTRIBUTED SYSTEMS: WORKSHOPS, SEVEN INTERNATIONAL  
CONFERENCE ON 04 July 2000 - 07 July 2000, LOS ALAMITOS, pages 307 -  
312;

ABSTRACT EP 1225530 A1

Numbers 1, 2, ..., M are assigned to bidding prices from the minimum to  
maximum values V1)) to VN)). For a bidding value Vvi)) each user 11-i  
generates two sequences of information si))=(si,1)), si,2, ...,)) si,M)))  
and ti))=(ti,1,)) ti,2,)) ..., ti,M))) such that si,1))=ti,1)), ...,  
si,vi-1))=ti,vi-1)), si,vi))(not equal to)ti,vi)), ..., si,M))(not equal  
to)ti,M)), then secretly sends the two sequences of information s; and t;  
to quantitative competition apparatuses 15A and 15B, respectively, and  
sends hash values H1i))=h(si)) and H2i))=h(ti)) of the two sequences of  
information si)) and ti)) and a hash value h(Vvi))(parallel to)ri)))  
containing an intended value Vvi)) to a bulletin board apparatus 21. The  
quantitative competition apparatuses 15A and 15B extract w-th elements  
si,w)) from respective sequences s1)) to sN)) and w-th elements ti,w))  
from respective sequences t1)) to tN)), then create a concatenation  
Seqs,w)) of N elements si,w)) and a concatenation Seqt,w)) of N elements  
ti,w)), then compare them using a one-way function without revealing  
their values, and, if they differ, deciding that the intended value Vvi))  
equal to or smaller than a value Vw)) is present, and determines the  
minimum value by changing w.

ABSTRACT WORD COUNT: 189

NOTE:

Figure number on first page: 1

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Application: 020724 A1 Published application with search report  
Examination: 020724 A1 Date of request for examination: 20020117  
Examination: 030416 A1 Date of dispatch of the first examination  
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Grant: 051026 B1 Granted patent  
Change: 061004 B1 Title of invention (German) changed: 20061004  
Change: 061004 B1 Title of invention (English) changed: 20061004  
Change: 061004 B1 Title of invention (French) changed: 20061004

LANGUAGE (Publication,Procedural,Application): English; English; English

FULLTEXT AVAILABILITY:

Available Text	Language	Update	Word Count
CLAIMS A	(English)	200230	5122
CLAIMS B	(English)	200543	4945
CLAIMS B	(German)	200543	4593
CLAIMS B	(French)	200543	5249
SPEC A	(English)	200230	9495
SPEC B	(English)	200543	9744
Total word count - document A			14620
Total word count - document B			24531
Total word count - documents A + B			39151

...SPECIFICATION to some extent, and if the number of bits of each element  
is large, the random numbers R1i)) and R2i)) may be omitted .

The quantitative competition apparatuses 15A (and 15B) comprise, as shown in Fig. 4 (and Fig. 5), receiving parts 40A (40B), decrypting parts 37A (47B), storage parts 41A (41B), random generating parts 42A (42B), hash function calculating...

...SPECIFICATION to some extent, and if the number of bits of each element is large, the random numbers R1i)) and R2i)) may be omitted.

The quantitative competition apparatuses 15A (and 15B) comprise, as shown in Fig. 4 (and Fig. 5), receiving parts 40A (40B), decrypting parts 37A (47B), storage parts 41A (41B), random generating parts 42A (42B), hash function calculating...

20/5,K/16 (Item 16 from file: 348)  
DIALOG(R)File 348:EUROPEAN PATENTS  
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01217622

Optical disk, optical disk recording and reproducing apparatus, and method for recording and reproducing

Optische Platte, optisches Plattenaufzeichnungs- und wiedergabegerat, und Verfahren zur Aufzeichnung und Wiedergabe

Disque optique, appareil pour l'enregistrement et la reproduction de disque optique, et methode d'enregistrement et de reproduction

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PATENT (CC, No, Kind, Date): EP 1058254 A2 001206 (Basic)  
EP 1058254 A3 011121  
EP 1058254 B1 040707

APPLICATION (CC, No, Date): EP 2000108910 000427;

PRIORITY (CC, No, Date): JP 99122104 990428; JP 99128197 990510; JP 99299635 991021

DESIGNATED STATES: DE; FR; GB

EXTENDED DESIGNATED STATES: AL; LT; LV; MK; RO; SI

RELATED DIVISIONAL NUMBER(S) - PN (AN):  
(EP 2004004964)

INTERNATIONAL PATENT CLASS (V7): G11B-020/00; G11B-007/007; G11B-023/28; G06F-001/00

CITED PATENTS (EP B): EP 442566 A; EP 802527 A; EP 954173 A; EP 984346 A;

WO /21087 A; WO 98/58368 A; GB 2332977 A; US 5513169 A; US 5596639 A; US 5646993 A; US 5745568 A; US 5752009 A  
CITED REFERENCES (EP B):  
PATENT ABSTRACTS OF JAPAN vol. 1997, no. 10, 31 October 1997 (1997-10-31)  
& JP 09 171619 A (SONY CORP), 30 June 1997 (1997-06-30)  
PATENT ABSTRACTS OF JAPAN vol. 2000, no. 07, 29 September 2000  
(2000-09-29) & JP 2000 113586 A (VICTOR CO OF JAPAN LTD), 21 April 2000  
(2000-04-21);

ABSTRACT EP 1058254 A2

An optical disk of recording type on which data is recordable is provided which includes a data recording and reproducing area (102) for recording data therein and reproducing data therefrom, and a read-only disk identification information area (104) for recording disk identification information (107) for identifying the optical disk therein. In the optical disk, the disk identification information (107) is formed by removing a reflection film formed on the optical disk in a strip shape. The disk identification information (107) includes an inherent disk identifier for each optical disk, and the data recording and reproducing area includes an area for recording therein encrypted data, which is encrypted using information including the disk identification information for identifying the optical disk as a key.

ABSTRACT WORD COUNT: 123

NOTE:

Figure number on first page: 1

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Application: 001206 A2 Published application without search report  
Change: 001220 A2 Title of invention (French) changed: 20001101  
Search Report: 011121 A3 Separate publication of the search report  
Examination: 020717 A2 Date of request for examination: 20020508  
Examination: 020925 A2 Date of dispatch of the first examination  
report: 20020807  
Change: 031022 A2 Title of invention (German) changed: 20030902  
Change: 031022 A2 Title of invention (English) changed: 20030902  
Change: 031022 A2 Title of invention (French) changed: 20030902  
Change: 040428 A2 Application number of divisional application  
(Article 76) changed: 20040309  
Grant: 040707 B1 Granted patent  
Change: 040804 B1 Inventor information changed: 20040611  
Change: 040804 B1 Inventor information changed: 20040611  
Oppn None: 050629 B1 No opposition filed: 20050408

LANGUAGE (Publication,Procedural,Application): English; English; English

FULLTEXT AVAILABILITY:

Available Text	Language	Update	Word Count
CLAIMS A	(English)	200049	5366
CLAIMS B	(English)	200428	1383
CLAIMS B	(German)	200428	1161
CLAIMS B	(French)	200428	1704
SPEC A	(English)	200049	30838
SPEC B	(English)	200428	28902
Total word count - document A			36210
Total word count - document B			33150
Total word count - documents A + B			69360

...SPECIFICATION so as to be divided into the first decipher key area 2201 having a division decipher key of 4 bytes and the second decipher key area 2202 having a division decipher key of 4 bytes. Therefore, in spite of the size of the encrypted content recorded...

...sectors, a plurality of sectors (2 sectors in Fig. 27) are utilized. In this case, dummy data is recorded in the unused area as complementary data. In an example of Fig. 27, complementary data 2203 for one...

...SPECIFICATION so as to be divided into the first decipher key area 2201 having a division decipher key of 4 bytes and the second decipher key area 2202 having a division decipher key of 4 bytes. Therefore, in spite of the size of the encrypted content recorded...

...sectors, a plurality of sectors (2 sectors in Fig. 27) are utilized. In this case, dummy data is recorded in the unused area as complementary data. In an example of Fig. 27, complementary data 2203 for one...

20/5,K/19 (Item 19 from file: 348)  
DIALOG(R)File 348:EUROPEAN PATENTS  
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01150840

Image data reproducing method, image data reproducing apparatus, image data recording method and image data recording apparatus  
Bilddatenwiedergabeverfahren, Bilddatenwiedergabeanlage, Bilddatenaufnahmeverfahren, und Bilddatenaufnahmeanlage  
Procede de reproduction de donnees d'image, appareil de reproduction de donnees d'image, procede d'enregistrement de donnees d'image, et appareil d'enregistrement de donnees d'image

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PATENT (CC, No, Kind, Date): EP 1003338 A2 000524 (Basic)  
EP 1003338 A3 020703

APPLICATION (CC, No, Date): EP 99122075 991116;

PRIORITY (CC, No, Date): JP 98326937 981117

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EXTENDED DESIGNATED STATES: AL; LT; LV; MK; RO; SI

INTERNATIONAL PATENT CLASS (V7): H04N-009/82

ABSTRACT EP 1003338 A2

A method of reproducing first images and second images simultaneously, synchronizing them with each other. The method includes the processes of: reading first image data representing the first images and second image data representing the second images from a recording medium (1); storing the first image data into a first memory device (58) and storing the second image data into a second memory device (59); and separately and simultaneously decoding the first image data and the second image data by using a first decoding device (61) and a second decoding device (62). In this method, each of the first image data and the second image data is divided into a plurality of data units (43) each having an equal time length and an equal data size, and each of the data units (43a) of the first image data and each of the data units (43b) of the second image data are alternately arranged on the recording medium (1).

ABSTRACT WORD COUNT: 159

NOTE:

Figure number on first page: 3

LEGAL STATUS (Type, Pub Date, Kind, Text):

Application: 000524 A2 Published application without search report  
Search Report: 020703 A3 Separate publication of the search report  
Examination: 020925 A2 Date of request for examination: 20020722  
Change: 070110 A2 Title of invention (German) changed: 20070110  
Change: 070110 A2 Title of invention (English) changed: 20070110  
Change: 070110 A2 Title of invention (French) changed: 20070110  
Change: 070801 A2 Title of invention (German) changed: 20070801  
Change: 070801 A2 Title of invention (English) changed: 20070801  
Change: 070801 A2 Title of invention (French) changed: 20070801

LANGUAGE (Publication,Procedural,Application): English; English; English

FULLTEXT AVAILABILITY:

Available Text	Language	Update	Word Count
CLAIMS A	(English)	200021	2065
SPEC A	(English)	200021	11293
Total word count - document A			13358
Total word count - document B			0
Total word count - documents A + B			13358

...SPECIFICATION rate, predetermined dummy data are written into this portion. At the time of reproduction, the dummy data are identified and are not used for the decoding process by means of MPEG2. In another way, the dummy data may be null data which are ignored at the time of the decoding process.

Image data are compressed and encoded so that a predetermined number of bits are...

20/5,K/21 (Item 21 from file: 348)

DIALOG(R)File 348:EUROPEAN PATENTS

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00943762

Method for the settlement of credit by an IC card

Verfahren zur Regelung von Gutschriften mittels einer Chipkarte

Methode d'etablissement de paiements par carte de credit

PATENT ASSIGNEE:

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Miyaguchi, Shoji, 1-4-4, Sugano,, Ichikawa-shi, Chiba, (JP)

Okamoto, Tatsuaki, 94-2-5-503, Nagasawa, Yokosuka-shi, Kanagawa, (JP)

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PATENT (CC, No, Kind, Date): EP 856822 A2 980805 (Basic)

EP 856822 A3 981021

EP 856822 B1 030305

APPLICATION (CC, No, Date): EP 98104504 930916;

PRIORITY (CC, No, Date): JP 92249293 920918; JP 92249294 920918; JP

92308688 921118; JP 92317254 921126; JP 92317255 921126

DESIGNATED STATES: DE; FR; GB

RELATED PARENT NUMBER(S) - PN (AN):

EP 588339 (EP 93114917)

INTERNATIONAL PATENT CLASS (V7): G07F-007/10

CITED PATENTS (EP B): EP 82958 A; EP 281059 A; EP 421409 A; EP 422230 A; EP



ABSTRACT EP 856822 A2

An IC card (6) has a card information memory area wherein there are written a card identification number IDU, a predetermined setting number Ns, a first master digital signature SA4 for said setting number Ns and a second master digital signature SA5 for information including said card identification number IDU and said first master digital signature SA4. An IC card terminal (2a, 2b) has terminal information memory area wherein there are written a master public key nA for verification of a master digital signature, terminal secret keys pT and qT for the creation of a terminal digital signature and a terminal public key nT for verification of said terminal digital signature. When inserted into the IC card terminal, said IC card transmits said card identification number IDU and said first and second master digital signatures to said IC card terminal; said IC card terminal verifies said second master digital signature SA5 and, if valid, instructs the entering of a password and transmits a password Nc' to said IC card when it is entered; said IC card matches said password Nc' with a password Nc stored in said card information memory means and, if they match, transmits an authentication signal to said IC card terminal; and upon receiving said authentication signal said IC card terminal becomes enabled for providing a service and, after completion of said service, records information including an amount value V for said service and said received card identification number IDU, as usage/management information in usage/management memory means.

ABSTRACT WORD COUNT: 251

NOTE:

Figure number on first page: 1

LEGAL STATUS (Type, Pub Date, Kind, Text):

Change: 020710 A2 Title of invention (German) changed: 20020523  
 Examination: 20000315 A2 Date of dispatch of the first examination report: 20000127  
 Oppn None: 040225 B1 No opposition filed: 20031208  
 Change: 020710 A2 Title of invention (French) changed: 20020523  
 Change: 020710 A2 Title of invention (English) changed: 20020523  
 Grant: 030305 B1 Granted patent  
 Application: 980805 A2 Published application (A1with Search Report ;A2without Search Report)  
 Examination: 980805 A2 Date of filing of request for examination: 980312  
 Change: 980909 A2 Inventor (change)  
 Change: 981007 A2 Inventor (change)  
 Search Report: 981021 A3 Separate publication of the European or International search report

LANGUAGE (Publication,Procedural,Application): English; English; English

FULLTEXT AVAILABILITY:

Available Text	Language	Update	Word Count
CLAIMS A	(English)	199832	882
CLAIMS B	(English)	200310	1106
CLAIMS B	(German)	200310	1008
CLAIMS B	(French)	200310	1351
SPEC A	(English)	199832	14787
SPEC B	(English)	200310	14923
Total word count - document A			15672
Total word count - document B			18388
Total word count - documents A + B			34060

...SPECIFICATION function to simplify the processing of the IC card system.

Also it is possible to omit either one of the random number R and X although security decreases. Conversely, by prestoring algorithms for encipherment of information to be transmitted and a common key for encipherment and decipherment in memories of the IC card 6 and the IC

card terminal 2, the mutual...

...SPECIFICATION function to simplify the processing of the IC card system.

Also it is possible to omit either one of the random number R and X although security decreases. Conversely, by prestoring algorithms for encipherment of information to be transmitted and a common key for encipherment and decipherment in memories of the IC card 6 and the IC card terminal 2, the mutual...

20/5,K/27 (Item 27 from file: 348)  
DIALOG(R)File 348:EUROPEAN PATENTS  
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00655909

Method and apparatus for waveform shaping of packet data  
Verfahren und Einrichtung zur Signalformung von Paketdaten  
Procede et dispositif de mise en forme de donnees en paquets  
PATENT ASSIGNEE:

MATSUSHITA ELECTRIC INDUSTRIAL CO., LTD., (216883), 1006, Oaza-Kadoma,  
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PATENT (CC, No, Kind, Date): EP 631398 A2 941228 (Basic)  
EP 631398 A3 000308  
EP 631398 B1 041117

APPLICATION (CC, No, Date): EP 94109760 940623;

PRIORITY (CC, No, Date): JP 93154776 930625; JP 93223292 930908; JP 9413760  
940207

DESIGNATED STATES: DE; FR; GB; NL

RELATED DIVISIONAL NUMBER(S) - PN (AN):

EP 1333633 (EP 2003008489)

EP 1331778 (EP 2003008498)

INTERNATIONAL PATENT CLASS (V7): H04L-025/03; H03G-003/30; H04L-027/20

CITED PATENTS (EP B): EP 369135 A; EP 463625 A; EP 494696 A; EP 545546 A;

EP 588424 A; DE 3941265 A; US 3912863 A; US 5140613 A; US 5412691 A

CITED REFERENCES (EP B):

PATENT ABSTRACTS OF JAPAN vol. 016, no. 475 (E-1273), 2 October 1992  
(1992-10-02) & JP 04 170128 A (MATSUSHITA ELECTRIC IND CO LTD), 17 June  
1992 (1992-06-17)

ALLAN R D ET AL: "A HIGH PERFORMANCE SATELLITE DATA MODEM USING REAL-TIME  
DIGITAL SIGNAL PROCESSING TECHNIQUES" JOURNAL OF THE INSTITUTION OF  
ELECTRONIC AND RADIO ENGINEERS, vol. 58, no. 3, 1 May 1988  
(1988-05-01), pages 117-124, XP000760573

PATENT ABSTRACTS OF JAPAN vol. 008, no. 261 (E-281), 30 November 1984  
(1984-11-30) & JP 59 132267 A (NIPPON DENKI KK), 30 July 1984  
(1984-07-30)

POLOZEC LE X ET AL: "A FAST AUTOMATIC GAIN CONTROL FOR A RF AMPLIFIER  
USED IN A DCS1800 (PCN) BASE STATION TRANSMITTER" PROCEEDINGS OF THE  
23RD. EUROPEAN MICROWAVE CONFERENCE, MADRID, SEPT. 6 - 9, 1993, 6  
September 1993 (1993-09-06), pages 935-938, XP000630037 EUROPEAN  
MICROWAVE CONFERENCE MANAGEMENT COMMITTEE ISBN: 0-946821-23-2

ABSTRACT EP 631398 A2

In the transmitter which carries out burst transmission using information data as a packet, if the status is divided into four modes, namely, burst stop mode, burst rising mode, burst continuous mode, and burst falling mode,

a waveform shaping equipment designed to read out shaped waveform data for each mode from outputs of either of the two memory tables, the first memory table which holds waveform data for specific data patterns used in common in burst rising mode and burst falling mode and the second memory table which holds waveform data for all data patterns used in the burst continuous mode, or a waveform shaping equipment comprising the third memory table which holds waveform data corresponding to all the data patterns used in the burst rising mode and the fourth memory table which holds waveform data corresponding to all data patterns used in the burst falling mode and generating shaped waveform data by synthesizing the two outputs of the third and the fourth memory tables at the time of burst continuous mode. (see image in original document)

ABSTRACT WORD COUNT: 179

NOTE:

Figure number on first page: 13

LEGAL STATUS (Type, Pub Date, Kind, Text):

Examination: 000802 A2 Date of request for examination: 20000607  
Change: 20000223 A2 International Patent Classification changed:  
20000106  
Oppn None: 051109 B1 No opposition filed: 20050818  
Change: 031029 A2 Inventor information changed: 20030906  
Examination: 021211 A2 Date of dispatch of the first examination  
report: 20021028  
Change: 020417 A2 International Patent Classification changed:  
20020301  
Change: 030604 A2 Application number of divisional application  
(Article 76) changed: 20030416  
Grant: 041117 B1 Granted patent  
Application: 941228 A2 Published application (A1with Search Report  
;A2without Search Report)  
Search Report: 20000308 A3 Separate publication of the search report  
Change: 990915 A2 International Patent Classification changed:  
19990729

LANGUAGE (Publication,Procedural,Application): English; English; English

FULLTEXT AVAILABILITY:

Available Text	Language	Update	Word Count
CLAIMS A	(English)	EPABF2	7821
CLAIMS B	(English)	200447	830
CLAIMS B	(German)	200447	699
CLAIMS B	(French)	200447	1020
SPEC A	(English)	EPABF2	19850
SPEC B	(English)	200447	7855
Total word count - document A			27676
Total word count - document B			10404
Total word count - documents A + B			38080

...SPECIFICATION transmitted. In this fourth embodiment, the hardware scale with respect to the shift register, address decoder, memory, etc. increases but no complicated control means is required to change over the memory...

...the information data into the shift register in the pattern generator from the first time without using dummy data, and the time equivalent to several time slots required for transmitting dummy data strings can...

...the same principle will be applied to multilevel patterns exceeding ternary patterns, if an address decoder for converting multilevel patterns to binary address signals is provided, partial waveforms of baseband signals...

20/5,K/31 (Item 31 from file: 348)  
DIALOG(R)File 348:EUROPEAN PATENTS  
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00361288

Decoder circuit.  
Dekodierschaltung.  
Circuit decodeur.

PATENT ASSIGNEE:

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PATENT (CC, No, Kind, Date): EP 327340 A2 890809 (Basic)  
EP 327340 A3 910320  
EP 327340 B1 931006

APPLICATION (CC, No, Date): EP 89300971 890201;

PRIORITY (CC, No, Date): JP 8822285 880202

DESIGNATED STATES: DE; FR; GB

INTERNATIONAL PATENT CLASS (V7): G11C-008/00;

CITED REFERENCES (EP A):

PATENT ABSTRACTS OF JAPAN, vol. 7, no. 78 (P-188) 1223 , 31st March 1983; & JP-A-58 6589 (HITACHI) 14-01-1983

PATENT ABSTRACTS OF JAPAN, vol. 10, no. 283 (E-440) 2339 , 26th September 1986; & JP-A-61 101 124 (HITACHI) 20-05-1986;

ABSTRACT EP 327340 A2

A decoder circuit comprises a plurality of signal conductors ( $X(\text{sub } 0) - X(\text{sub } 7)$ , DL), first potential setting units (Tpc) for setting a potential of the signal conductors ( $X(\text{sub } 0) - X(\text{sub } 7)$ , DL) to a first potential, second potential setting units (Cs) for maintaining a potential of one signal conductor ( $X(\text{sub } 0) - X(\text{sub } 7)$ ) at the first potential and bringing a potential of the remaining signal conductors ( $X(\text{sub } 0) - X(\text{sub } 7)$ , DL) to a second potential, and a respective transfer transistor ( $T(\text{sub } 0) - T(\text{sub } 7)$ ) provided for each signal conductor ( $X(\text{sub } 0) - X(\text{sub } 7)$ ). A gate electrode of each transfer transistor ( $T(\text{sub } 0) - T(\text{sub } 7)$ ) is connected to its respective signal conductor ( $X(\text{sub } 0) - X(\text{sub } 7)$ ), a source electrode thereof is connected to another signal conductor ( $X(\text{sub } 1) - X(\text{sub } 7)$ , DL), and one transfer transistor ( $T(\text{sub } 0) - T(\text{sub } 7)$ ) is switched ON and a decoded output signal is output when a gate potential of that transfer transistor ( $T(\text{sub } 0) - T(\text{sub } 7)$ ) is maintained at the first potential and a source potential thereof is brought to the second potential. Therefore, only one output signal conductor ( $X(\text{sub } 0) - X(\text{sub } 7)$ ) is selected during a decoder operation without using a discharge detecting circuit of a nonselected output signal conductor and a timing signal, and thus an operation speed of the decoder circuit is much higher and the decoding circuit has a simpler construction and is suitable as a large scale integration circuit.

ABSTRACT WORD COUNT: 258

LEGAL STATUS (Type, Pub Date, Kind, Text):

Application: 890809 A2 Published application (A1with Search Report ;A2without Search Report)

Search Report: 910320 A3 Separate publication of the European or  
International search report  
Examination: 910717 A2 Date of filing of request for examination:  
910516  
Examination: 930310 A2 Date of despatch of first examination report:  
930126  
Grant: 931006 B1 Granted patent  
Oppn None: 940928 B1 No opposition filed  
LANGUAGE (Publication,Procedural,Application): English; English; English  
FULLTEXT AVAILABILITY:  
Available Text Language Update Word Count  
CLAIMS B (English) EPBBF1 1313  
CLAIMS B (German) EPBBF1 964  
CLAIMS B (French) EPBBF1 1207  
SPEC B (English) EPBBF1 5528  
Total word count - document A 0  
Total word count - document B 9012  
Total word count - documents A + B 9012

...SPECIFICATION level.

Figures 6 and 7 are circuit diagrams showing modifications of the  
embodiment of the decoder circuit shown in Fig. 5. In these decoder  
circuits, the dummy signal conductor DL can be omitted.

As shown in Fig. 6, in one modification of the decoder circuit,  
output signal conductors X(sub 0) - X(sub 7) are divided into pairs of  
two adjacent...

...first output signal conductor X(sub 0).

As described above, in these modifications of the decoder circuit,  
the dummy signal conductor DL can be omitted. Furthermore, a  
source electrode of each of the transfer transistors T(sub 0) - T(sub 7)  
is...

20/5,K/36 (Item 36 from file: 349)  
DIALOG(R)File 349:PCT FULLTEXT  
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01157716 \*\*Image available\*\*

SYSTEM AND METHOD FOR DATA ENCRYPTION  
SYSTEME ET PROCEDE DE CRYPTAGE DE DONNEES

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Patent and Priority Information (Country, Number, Date):

Patent: WO 200479980 A2-A3 20040916 (WO 0479980)

Application: WO 2004US6688 20040305 (PCT/WO US04006688)

Priority Application: US 2003384147 20030305

Designated States:

(All protection types applied unless otherwise stated - for applications  
2004+)

AE AG AL AM AT AU AZ BA BB BG BR BW BY BZ CA CH CN CO CR CU CZ DE DK DM  
DZ EC EE EG ES FI GB GD GE GH GM HR HU ID IL IN IS JP KE KG KP KR KZ LC  
LK LR LS LT LU LV MA MD MG MK MN MW MX MZ NA NI NO NZ OM PG PH PL PT RO  
RU SC SD SE SG SK SL SY TJ TM TN TR TT TZ UA UG US UZ VC VN YU ZA ZM ZW

(EP) AT BE BG CH CY CZ DE DK EE ES FI FR GB GR HU IE IT LU MC NL PL PT RO  
SE SI SK TR

(OA) BF BJ CF CG CI CM GA GN GQ GW ML MR NE SN TD TG

(AP) BW GH GM KE LS MW MZ SD SL SZ TZ UG ZM ZW

(EA) AM AZ BY KG KZ MD RU TJ TM

Main International Patent Class (v7): H04N-007/167

Publication Language: English

Filing Language: English

Fulltext Availability:

Detailed Description

Claims

Fulltext Word Count: 12794

#### English Abstract

An encryption technology uses an encryption map and video graphics processing technology to combine an encryption map with the image data. Video image processing using maps, such as texture maps, bump maps and the like are combined with an encryption map to generate an encrypted image data. The image is subsequently decrypted using encryption keys that are combined with the encrypted image data using video graphics processing technologies. The decryption map is combined with the encrypted signal to generate a viewable image. The encryption keys may be on a per pixel basis with a separate encryption key for each pixel. Alternatively, an encryption key may contain active keys that may be used to decrypt more than one pixel as well as decoy keys that do not decrypt the signal and thus confound unauthorized decryption attempts. Using encryption maps, single video frames or portions of video frames may be encrypted. The encryption maps may be used with scaled image data and with two dimensional or three dimensional graphics processors.

#### French Abstract

L'invention porte sur une technique de cryptage utilisant une carte de cryptage et une technique de traitement de graphiques video. Le traitement des images video recourt a des cartes de texture, de reliefs, etc, combinees avec une carte de cryptage de maniere a obtenir des donnees d'image cryptees. L'image est ensuite decryptee a l'aide de clefs de cryptage qui sont combinees avec les donnees d'image cryptees a l'aide de techniques video de traitement graphique. La carte de decryptage est combinee au signal crypte pour donner une image visible. Il peut y avoir une clef de cryptage separee par pixel. En variante, une clef de cryptage peut contenir des clefs actives pouvant decrypter plus d'un pixel et des clefs leurres qui ne decryptent pas le signal et dejouent les tentatives de decryptage non autorisees. On peut a l'aide de cartes de cryptage crypter des images video uniques ou des parties d'images video. Les cartes de cryptage peuvent etre utilisees avec des donnees d'images agrandies ou reduites et avec des processeurs graphiques en 2D ou en 3D.

#### Legal Status (Type, Date, Text)

Publication 20040916 A2 Without international search report and to be republished upon receipt of that report.

Search Rpt 20041021 Late publication of international search report

Republication 20041021 A3 With international search report.

Republication 20041021 A3 Before the expiration of the time limit for amending the claims and to be republished in the event of the receipt of amendments.

Search Rpt 20041021 Late publication of international search report

Claim Mod 20041223 Later publication of amended claims under Article 19 received: 20041104

Republication 20041223 A3 With international search report.

Republication 20041223 A3 With amended claims.

Fulltext Availability:

Claims

Claim

... claim 12 wherein the encrypted output image data comprises a plurality of pixels and the decryption map has a plurality of storage locations corresponding to each of the plurality of pixels in the encrypted output image data and contains active decryption data in at least one location in the decryption map, the remaining locations in the decryption map containing decoy decryption data not used to decrypt the encrypted output image data.

18 The system of claim 17 wherein the active decryption data in the at least one location is stored in a predetermined one of the plurality of locations within the decryption map.

19 The system of claim 17 wherein the active decryption data in the at least one location is stored in a pseudo-random one of...

20/5,K/38 (Item 38 from file: 349)  
DIALOG(R)File 349:PCT FULLTEXT  
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01126380 \*\*Image available\*\*

METHOD AND SYSTEM FOR FORWARDING A CONTROL INFORMATION

PROCEDE ET SYSTEME DE TRANSFERT D'UNE INFORMATION DE COMMANDE

Patent Applicant/Assignee:

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FI (Nationality)

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MALKAMAKI Esa, Riippakoivuntie 17 B, FIN-02130 Espoo, FI,

Legal Representative:

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Munich, DE,

Patent and Priority Information (Country, Number, Date):

Patent: WO 200449749 A2-A3 20040610 (WO 0449749)

Application: WO 2003IB5361 20031112 (PCT/WO IB03005361)

Priority Application: US 2002304949 20021127

Designated States:

(Protection type is "patent" unless otherwise stated - for applications prior to 2004)

AE AG AL AM AT AU AZ BA BB BG BR BY BZ CA CH CN CO CR CU CZ DE DK DM DZ  
EC EE EG ES FI GB GD GE GH GM HR HU ID IL IN IS JP KE KG KP KR KZ LC LK  
LR LS LT LU LV MA MD MG MK MN MW MX MZ NI NO NZ OM PG PH PL PT RO RU SC  
SD SE SG SK SL SY TJ TM TN TR TT TZ UA UG UZ VC VN YU ZA ZM ZW

(EP) AT BE BG CH CY CZ DE DK EE ES FI FR GB GR HU IE IT LU MC NL PT RO SE  
SI SK TR

(OA) BF BJ CF CG CI CM GA GN GQ GW ML MR NE SN TD TG

(AP) BW GH GM KE LS MW MZ SD SL SZ TZ UG ZM ZW

(EA) AM AZ BY KG KZ MD RU TJ TM

Main International Patent Class (v7): H04Q-007/38

Publication Language: English

Filing Language: English

Fulltext Availability:

Detailed Description

Claims

Fulltext Word Count: 8163

English Abstract

The present invention relates to a method and system for forwarding a control information in a transmission signal of a communication network. A dummy information is provided in at least one predetermined portion of the transmission signal, and is replaced at least partly by the control information at a control device arranged on the transmission path of the transmission signal. Signaling space is thus generated by creating the

dummy information. Thereby, a fast control signaling can be provided which does not have to be originated at a network controlling functionality. Furthermore, if a dedicated link is used, less power is required and designing of new physical channel types is not required.

#### French Abstract

La presente invention se rapporte a un procede et un systeme de transfert d'une information de commande dans un signal de transmission d'un reseau de communication. Une information factice est acheminee dans au moins une partie predeterminee du signal de transmission et est remplacee au moins partiellement par l'information de commande au niveau d'un dispositif de commande dispose sur le trajet de transmission. Un espace de signalisation est alors genere par la creation de l'information factice. En consequence, une signalisation de commande rapide peut etre acheminee et il n'est pas necessaire qu'elle soit creee pour une fonctionnalite de commande de reseau. Par ailleurs, si la liaison utilisee est specialisee, on necessite moins de puissance et il n'est pas utile de designer de nouveaux types de canaux physiques.

#### Legal Status (Type, Date, Text)

Publication 20040610 A2 without international search report and to be republished upon receipt of that report.

Search Rpt 20040902 Late publication of international search report

Republication 20040902 A3 with international search report.

Examination 20041014 Request for preliminary examination prior to end of 19th month from priority date

#### Fulltext Availability: Detailed Description

#### Detailed Description ... slot.

Thus, in case of the fixed positions, portions which are most likely to be unused can be used as dummy data positions for forwarding the specific control signaling.

Due to the fixed positions, the bits are...

...a later interleaving function will distribute the bits. The use of flexible positions generally requires decoding of the transmission format indication information. However, if the dummy information is inserted to the...

...1 0, the whole frame F has to be received before this signaling can be decoded.

If the signaling bits are transmitted in selected time slots, the dummy bits D are...

...e. the second interleaving in the present WCDMA system, when the final positions are known. Non-used or non-replaced dummy bits D may preferably be replaced by IDTX indication bits, i.e., they are not transmitted...

20/5,K/39 (Item 39 from file: 349)  
DIALOG(R)File 349:PCT FULLTEXT  
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01028558 \*\*Image available\*\*

PURCHASING AID LOGISTICS APPLIANCE AND METHOD TO USE SAME  
INSTRUMENT LOGISTIQUE D'AIDE A L'ACHAT ET PROCEDE D'UTILISATION ASSOCIE  
Patent Applicant/Assignee:

LOCKHEED MARTIN CORPORATION, c/o Lockheed Martin Federal Systems, 1801



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Patent and Priority Information (Country, Number, Date):

Patent: WO 200358529 A1 20030717 (WO 0358529)  
Application: WO 2002US41559 20021227 (PCT/WO US0241559)  
Priority Application: US 200237382 20020104

Designated States:

(Protection type is "patent" unless otherwise stated - for applications  
prior to 2004)

AE AG AL AM AT AU AZ BA BB BG BR BY BZ CA CH CN CO CR CU CZ DE DK DM DZ  
EC EE ES FI GB GD GE GH GM HR HU ID IL IN IS JP KE KG KP KR KZ LC LK LR  
LS LT LU LV MA MD MG MK MN MW MX MZ NO NZ OM PH PL PT RO RU SC SD SE SG  
SK SL TJ TM TN TR TT TZ UA UG UZ VC VN YU ZA ZM ZW

(EP) AT BE BG CH CY CZ DE DK EE ES FI FR GB GR IE IT LU MC NL PT SE SI SK  
TR

(OA) BF BJ CF CG CI CM GA GN GQ GW ML MR NE SN TD TG

(AP) GH GM KE LS MW MZ SD SL SZ TZ UG ZM ZW

(EA) AM AZ BY KG KZ MD RU TJ TM

Main International Patent Class (v7): G06F-017/60

Publication Language: English

Filing Language: English

Fulltext Availability:

Detailed Description

Claims

Fulltext Word Count: 10848

English Abstract

A purchasing aid logistics appliance (10) that assists a purchaser with a shopping list generation, in-store product location, automated checkout, and financial management. The purchasing aid logistics appliance (10) is in the form of a battery (100) operated purchasing aid logistics appliance with touch screen display (50) and multimedia input/output. The purchasing aid logistics appliance (10) functions as a fiduciary aid and two-way communications device secure for managing money and processing perishable data. The purchasing aid logistics appliance (10) can function as a stand-alone device or synchronized with a user's personal computer (68) via a RF or IR link.

French Abstract

L'invention concerne un instrument logistique d'aide a l'achat (10) qui sert a un acheteur pour etablir une liste de provisions, localiser un produit en magasin, effectuer un controle de sortie automatise et realiser la gestion financiere. Cet instrument logistique d'aide a l'achat (10) se presente sous forme d'un dispositif a pile (100), dote d'un ecran tactile (50) et d'une entree/sortie multimedia. Ledit instrument logistique d'aide a l'achat (10) fonctionne comme dispositif d'aide fiduciaire et de dialogue securise pour la gestion financiere et le traitement de donnees temporaires. L'instrument logistique d'aide a l'achat (10) selon l'invention peut fonctionner en tant que dispositif autonome ou synchronise avec l'ordinateur personnel (68) d'un utilisateur par l'intermediaire d'une liaison RF ou IR.

Legal Status (Type, Date, Text)

Publication 20030717 A1 with international search report.

Patent and Priority Information (Country, Number, Date):

Patent: ... 20030717

Fulltext Availability:  
Detailed Description  
Publication Year: 2003

Detailed Description

... signals 50 but is instead computed as select lines for B010 (B00I+F) by address decoder 55. The central processor 11 thinks it is writing B000 through B003 four successive locations the address decoder 55 now gets a signal not to use the random number encode key 52 generated by random number generation circuitry 51. Now the central processor 11...

...commands by memory bus address signals 50 beginning with location B000 through B003. The address decoder 55 functions normally but the data retrieved from these four locations B000-B003 is not...

20/5,K/50 (Item 50 from file: 349)  
DIALOG(R)File 349:PCT FULLTEXT  
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00753871 \*\*Image available\*\*

OPTICAL DISK, OPTICAL DISK RECORDING AND REPRODUCING APPARATUS, METHOD FOR RECORDING, REPRODUCING AND DELETING DATA ON OPTICAL DISK, AND INFORMATION PROCESSING SYSTEM

DISQUE OPTIQUE, APPAREIL D'ENREGISTREMENT ET DE REPRODUCTION SUR DISQUE OPTIQUE, PROCEDE D'ENREGISTREMENT, REPRODUCTION ET EFFACEMENT DE DONNEES SUR DISQUE OPTIQUE, ET SYSTEME DE TRAITEMENT D'INFORMATIONS

Patent Applicant/Assignee:

MATSUSHITA ELECTRIC INDUSTRIAL CO LTD, 1006, Oaza Kadoma, Kadoma-shi, Osaka 571-8501, JP, JP (Residence), JP (Nationality)

Inventor(s):

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NAKATA Kouhei, B305, Eminensu Marunouchi, 13-15, Marunouchi-cho, Kawanishi-shi, Hyogo 666-0003, JP,

Legal Representative:

AOYAMA Tamotsu (et al) (agent), Aoyama & Partners, IMP Building, 3-7, Shiromi 1-chome, Chuo-ku, Osaka-shi, Osaka 540-0001, JP,

Patent and Priority Information (Country, Number, Date):

Patent: WO 200067257 A2-A3 20001109 (WO 0067257)

Application: WO 2000JP2750 20000427 (PCT/WO JP0002750)

Priority Application: JP 99122104 19990428; JP 99128197 19990510; JP 99299635 19991021

Designated States:

(Protection type is "patent" unless otherwise stated - for applications

prior to 2004)

CN KR

Main International Patent Class (v7): G11B-020/00

International Patent Class (v7): G11B-007/007; G11B-023/28; G06F-001/00

Publication Language: English

Filing Language: English

Fulltext Availability:

Detailed Description

Claims

Fulltext Word Count: 40940

#### English Abstract

An optical disk of recording type on which data is recordable is provided which includes a data recording and reproducing area (102) for recording data therein and reproducing data therefrom, and a read-only disk identification information area (104) for recording disk identification information (107) for identifying the optical disk therein. In the optical disk, the disk identification information (107) is formed by removing a reflection film formed on the optical disk in a strip shape. The disk identification information (107) includes an inherent disk identifier for each optical disk, and the data recording and reproducing area includes an area for recording therein encrypted data, which is encrypted using information including the disk identification information for identifying the optical disk as a key.

#### French Abstract

La presente invention concerne un disque optique de type a enregistrement sur lequel on peut enregistrer des donnees. Ce disque optique comprend une zone d'enregistrement et de reproduction de donnees concue pour y enregistrer des donnees et reproduire des donnees a partir de celle-ci, et une zone d'informations d'identification de disque ROM concue pour enregistrer les informations d'identification de disque permettant d'identifier le disque optique en question. Dans le disque optique, on constitue les informations d'identification de disque en retirant un film reflechissant applique sur le disque optique sous forme de bande. Les informations d'identification de disque comprennent un identificateur de disque inherent pour chaque disque optique, la zone d'enregistrement et de reproduction de donnees comprenant une zone d'enregistrement des donnees chiffrees, ces donnees ayant ete chiffrees en utilisant ces informations, y compris les informations d'identification de disque permettant d'identifier le disque optique comme une cle.

#### Legal Status (Type, Date, Text)

Publication 20001109 A2 without international search report and to be republished upon receipt of that report.

Examination 20001228 Request for preliminary examination prior to end of 19th month from priority date

Search Rpt 20010426 Late publication of international search report

Republication 20010426 A3 with international search report.

#### Patent and Priority Information (Country, Number, Date):

Patent: ... 20001109

Fulltext Availability:

Detailed Description

Publication Year: 2000

#### Detailed Description

... so

as to be divided into the first decipher key area 2201 having a division decipher key of 4 bytes and the second decipher key area 2202 having a division decipher key of 4 bytes. Therefore, in spite of the size of the encrypted content recorded...

...sectors, a plurality of sectors (2 sectors in Fig. 27) are utilized. In this case, dummy data is recorded in the unused

area as complementary data. In an example of Fig. 27t  
complementary data 2203 for one...

20/5,K/51 (Item 51 from file: 349)  
DIALOG(R)File 349:PCT FULLTEXT  
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00749017 \*\*Image available\*\*

METHODS AND APPLIANCES FOR ENCRYPTION SYSTEM VARYING DYNAMICALLY DEPENDING  
UPON VARIABLES AND ITS APPLICATIONS  
PROCEDES ET DISPOSITIFS POUR SYSTEME DE CHIFFREMENT A VARIATION DYNAMIQUE  
DEPENDANT DE VARIABLES ET LEURS APPLICATIONS

Patent Applicant/Inventor:

YU Choonyeol, 712-1201 Uruk Apt., Goongrae-dong, Goonpo-si, Gyeonggi-Do  
435-047, KR, KR (Residence), KR (Nationality)

Patent and Priority Information (Country, Number, Date):

Patent: WO 200062458 A2 20001019 (WO 0062458)

Application: WO 2000KR363 20000414 (PCT/WO KR0000363)

Priority Application: KR 9913182 19990414

Designated States:

(Protection type is "patent" unless otherwise stated - for applications  
prior to 2004)

AE AG AL AM AT AU AZ BA BB BG BR BY CA CH CN CR CU CZ DE DK DM DZ EE ES  
FI GB GD GE GH GM HR HU ID IL IN IS JP KE KG KP KR KZ LC LK LR LS LT LU  
LV MA MD MG MK MN MW MX NO NZ PL PT RO RU SD SE SG SI SK SL TJ TM TR TT  
TZ UA UG US UZ VN YU ZA ZW

(EP) AT BE CH CY DE DK ES FI FR GB GR IE IT LU MC NL PT SE

(OA) BF BJ CF CG CI CM GA GN GW ML MR NE SN TD TG

(AP) GH GM KE LS MW SD SL SZ TZ UG ZW

(EA) AM AZ BY KG KZ MD RU TJ TM

Main International Patent Class (v7): H04L

Publication Language: English

Filing Language: Korean

Fulltext Availability:

Detailed Description

Claims

Fulltext Word Count: 4516

English Abstract

The frequencies of commercial transactions or the kind of activities on a network system have been increased gradually such as transaction with credit cards, connecting onto a banking computer system, usage of vending machines, etc. For this purpose, this invention is a password algorithm system utilizing regularly varying elements such as hour, date, week day, month, etc. that change continuously and other kinds of elements so that the password changes according to the variable elements of time and other elements. In addition, by implementing dynamic variable elements being linked to user's ID, this invention disables the abuse of private information caught by someone on the Internet as ID. Someone who gets the data as ID but not its algorithm couldn't use it as it's not accepted in a server computer system.

French Abstract

La fréquence des transactions commerciales ou activités analogues sur des systèmes en réseau est en progression constante, telles que les transactions avec cartes de crédit, connexions à des systèmes informatiques bancaires, utilisation de distributeurs automatiques, etc. A cet effet, l'invention concerne un système d'algorithme de mot de passe utilisant régulièrement des éléments variables tels que l'heure, la date, le jour de la semaine, le mois, etc. qui changent continuellement, ainsi que d'autres types d'éléments, de sorte que le mot de passe change en fonction des éléments variables de temps et des autres éléments. De plus, par la mise en application d'éléments variables dynamiques liés à

l'identite de l'utilisateur, l'invention permet d'eliminer l'utilisation abusive d'informations privees, telle l'identite, pouvant etre interceptees par une personne sur Internet. Une personne qui intercepterait les donnees d'identite mais pas leur algorithme ne pourraient pas les utiliser, car le systeme d'ordinateur de serveur ne les accepte pas.

Legal Status (Type, Date, Text)

Publication 20001019 A2 In English translation (filed in Korean).  
Publication 20001019 A2 Without international search report and to be republished upon receipt of that report.  
Examination 20010111 Request for preliminary examination prior to end of 19th month from priority date

Patent and Priority Information (Country, Number, Date):

Patent: ... 20001019

Fulltext Availability:

Detailed Description  
Publication Year: 2000

Detailed Description

... the fingerprints and compare it to the fingerprint data that is saved.  
This variable coding/ decoding system has such features.

[ Clause 6 1

As for clause 1, password-complicating unit is included additionally with the other password elements mentioned to intensify security. This variable coding/ decoding system has such features.

[ Clause 7 1

As for clause 6, the password-complicating unit includes a dummy digit that is not used by the computer when decoding . This variable coding/ decoding system has such features.

[ Clause 8 1

It is made up of an input unit...

...be combined using arithmetical logic units. This is the characteristic of the system for coding/ decoding of the data.

[ Clause 15

As for clause 12, password-complicating unit is included additionally...

...password elements mentioned to intensify security. This is the characteristic of the system for coding/ decoding of the data.

[ Clause 16 1

As for clause 15, the password-complicating unit includes a dummy digit that is not used by the computer when decoding . This is the characteristic of the system for coding/ decoding of the data.

[ Clause 17 1

A variable password lock device constituted of CPU, data...

...intensify security.

[ Clause 23 1

As for clause 22, the password-complicating unit includes a dummy digit that is not used by the computer when decoding .

[ Clause 24 1

A data coding/ decoding device that is made up of CPU and a data coding/ decoding unit, codes/ decodes data based upon variable elements that are created naturally or artificially.

Clause 25 1

As for clause 24, the data coding/ decoding unit is saved in a data-storing device and is included additionally in the data...

...intensify security.

[ Clause 28 1

As for clause 27, the password-complicating unit includes a dummy digit that is not used by the computer when decoding .

[ Clause 29 1

A fingerprint recognizing system designed

20/5,K/61 (Item 61 from file: 349)

DIALOG(R)File 349:PCT FULLTEXT

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00120560 \*\*Image available\*\*

COMPATIBLE DESCRAMBLER FOR TELEVISION SYNCHRONIZATION

DEBROUILLEUR COMPATIBLE POUR SYNCHRONISATION DE TELEVISION

Patent Applicant/Assignee:

DeRIDDER Dennis W,

MILLER Donald J,

Inventor(s):

DeRIDDER Dennis W,

MILLER Donald J,

Patent and Priority Information (Country, Number, Date):

Patent: WO 8403811 A1 19840927

Application: WO 84US402 19840316 (PCT/WO US8400402)

Priority Application: US 83961 19830316

Designated States:

(Protection type is "patent" unless otherwise stated - for applications prior to 2004)

GB JP

Main International Patent Class (v7): H04N-007/16

Publication Language: English

Fulltext Availability:

Detailed Description

Claims

Fulltext Word Count: 10119

English Abstract

A descrambling circuit for video signals which have been rendered secure against unauthorized use by suppression or elimination of the horizontal sync signals generates (29) a decoding signal at the receiving end by detection (18, 20) of an unmodified portion of the video signal occurring during the vertical blanking interval or by detection (40) of the 3.58 MHz color burst signal. In this way, descrambling can be accomplished without need for reference to a decoding signal with the scrambled video or a key inserted in the video signal at the transmission end. This also permits a single descrambling circuit to decode video signals which have been scrambled by different techniques.

French Abstract

Un circuit de debrouillage pour signaux video qui ont ete premunis contre toute utilisation non autorisee par la suppression ou l'elimination des signaux de synchronisation horizontale emet (29) un signal de decodage a l'extremite receptrice lors de la detection (18, 20)

d'une partie non modifiée du signal video se presentant pendant l'intervalle de suppression verticale ou lors de la detection (40) du signal de synchronisation de couleurs de 3,58 MHz. De cette maniere, il est possible d'effectuer un debrouillage sans avoir besoin de se referer a un signal de decodage lors d'un brouillage video ou de l'insertion d'une clef dans le signal video a l'extremite de transmission. Cela permet egalement a un seul circuit de debrouillage de decoder des signaux video qui ont ete brouilles par des procedes differents.

Patent and Priority Information (Country, Number, Date):

Patent: ... 19840927

Fulltext Availability:

Claims

Publication Year: 1984

Claim

... of the horizontal sync pulses is prevented; and  
transmitting said scrambled video signal-and a- false  
decoding signal which cannot be used to accurately descramble  
the scrambled video signal.

53 A method according to claim 52, further comprising  
the-steps...

...unmodified signal portion of said  
scrambled video signal not forming a part of said false  
decoding signal;  
generating a decoding signal solely on the basis of  
said detected unmodified signal portion; and  
combining said scrambled...

?

File 347:JAPIO Dec 1976-2007/Jun(Updated 070926)

(c) 2007 JPO & JAPIO

File 350:Derwent WPIX 1963-2007/UD=200764

(c) 2007 The Thomson Corporation

Set	Items	Description
S1	343327	FALSE OR INVALID OR NULL OR PSUEDO OR PSEUDO OR RANDOM OR - RAND OR CHAFF OR SEMIRANDOM OR PSUEDORAND? OR PSEUDORAND? OR - DUMMY OR NONCE? ? OR DECOY?
S2	2540	SHILL?? OR SPOOF?? OR PHONEY? ? OR PHONY? ? OR PHONIE? ? - OR FAKE? ? OR SHAM OR SHAMS
S3	30597	S1:S2(1W)(VALUE OR VALUES OR NUMBER? ? OR NUMERAL? ? OR NU- MERIC?? OR ALPHANUMERIC? OR QUANTITY? ? OR QUANTITIES OR CHAR- ACTER? ? OR BIT OR BITS OR DATA OR DATUM? ?)
S4	25500	S1:S2(1W)(DIGIT? ? OR SIGNAL? ? OR PULSE OR PULSES OR INTE- GER? ? OR STRING OR STRINGS OR SUBSTRING? ? OR SEQUENCE OR SE- QUENCES OR SUBSEQUENCE? ? OR NOISE OR VARIABLE? ?)
S5	13080923	USED OR USE OR USING OR USAGE? ? OR EMPLOY? OR UTILIS??? OR UTILISATION? ? OR UTILIZ??? OR UTILIZATION? ?
S6	492	NONEMPLOY? OR UNEMPLOY? OR NONUTILIS? OR NONUTILIZ? OR UNU- TILIS? OR UNUTILIZ?
S7	25372	NONUSE? OR NONUSING OR NONUSAGE? OR UNUSE? OR UNUSING OR U- NUSAGE? OR DISUSE? ? OR DISUSING OR DISUSAGE?
S8	330160	('NOT' OR WITHOUT OR NEVER OR CANNOT OR T) (1W)S5
S9	44477	(NON OR UN OR DIS)()S5
S10	133	S3:S4(5N)(OMIT? OR OMISSION? ? OR EXCLUD? OR EXCLUSION?)
S11	193	S3:S4(5N)S8:S9
S12	220709	DECRYPT? OR UNCRYPT? OR UNENCRYPT? OR DECOD??? OR UNENCOD?- ?? OR UNCOD??? OR UNENCRYPT? OR UNENCIPHER? OR UNENCYIPHER? OR DECIPHER? OR DECYPHER?
S13	39938	UNSCRAMBL? OR DESCramBL? OR UNCIPHER? OR UNCYPHER? OR DECO- MPRESS? OR DEPACK? OR UNPACK? OR UNCOMPRESS? OR DECOMPACT? OR UNCOMPACT?
S14	28	S3:S4(5N)S6:S7
S15	49	(S10:S11 OR S14) AND S12:S13
S16	34	S15 AND PY=1963:2003
S17	20	S15 AND AY=1963:2003 AND AC=US
S18	37	S16:S17

? t18/9/2,4,10-12

18/9/2 (Item 2 from file: 347)

DIALOG(R)File 347:JAPIO

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06204393 \*\*Image available\*\*

DIGITAL PRIVACY TELEPHONE SET

PUB. NO.: 11-145950 [JP 11145950 A]

PUBLISHED: May 28, 1999 ( 19990528)

INVENTOR(s): KODAMA HIROHISA

APPLICANT(s): NEC CORP

APPL. NO.: 09-312463 [JP 97312463]

FILED: November 13, 1997 (19971113)

INTL CLASS: H04L-009/20; H04L-009/12; H04L-009/36

#### ABSTRACT

PROBLEM TO BE SOLVED: To provide a digital privacy telephone set which can reduce speaking faults caused by the omission or erroneous insertion of data until transmission data from a caller are received on the side of a callee during privacy communication.

SOLUTION: On the side of the caller, phase information is added to



enciphered data, on the side of the callee, the phase information is detected and by comparing plural pieces of continuously or intermittently received phase information, the presence/absence of the omission or erroneous insertion of received data is detected and corrected. Besides, restored audio data outputted from a deciphering part 33 are received as input signals, and a data converting part 39 is provided for converting a data correspondent part and a phase information correspondent part, for which the omission is detected by data phase correcting means 37 and 38, to soundless data and outputting them. Thus, dummy data are inserted to the omission detected enciphered data part and after descramble processing is performed to that dummy data part and the phase information part, they are converted to the soundless data by the data converting means.

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18/9/4 (Item 4 from file: 347)  
DIALOG(R)File 347:JAPIO  
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05730306 \*\*Image available\*\*  
DIGITAL SECRECY DEVICE

PUB. NO.: 10-013406 [JP 10013406 A]  
PUBLISHED: January 16, 1998 ( 19980116)  
INVENTOR(s): KODAMA HIROHISA  
APPLICANT(s): NEC CORP [000423] (A Japanese Company or Corporation), JP  
(Japan)  
APPL. NO.: 08-163962 [JP 96163962]  
FILED: June 25, 1996 (19960625)  
INTL CLASS: [6] H04L-009/36; H04L-009/20; H04M-001/68  
JAPIO CLASS: 44.3 (COMMUNICATION -- Telegraphy); 44.2 (COMMUNICATION --  
Transmission Systems); 44.4 (COMMUNICATION -- Telephone)

#### ABSTRACT

PROBLEM TO BE SOLVED: To hold the synchronism of pseudo random numbers in a scramble means and a descramble part by adding phase information to scramble data in a transmitter-side, detecting phase information on a receiver-side and comparing a plurality of pieces of phase information which are continuously received.

SOLUTION: Phase information C generated in a data phase control part 26 is added to scramble data B from the scramble part 22 and it is transmitted to a transmission data control part 15. On the receiver-side, the presence or absence of the omission and the erroneous insertion of scramble data E is detected in a data phase detection part 37, based on phase information separated in a received data control part 36. A data phase correction part 38 estimates data quantity which is omitted and erroneously inserted by a phase correction signal from the data phase detection part 37, inserts dummy data in the omitted period of scramble data E and eliminates the data part of erroneously inserted scramble data E. A data conversion part 39 converts dummy data into silence data.

18/9/10 (Item 10 from file: 347)  
DIALOG(R)File 347:JAPIO  
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02916268 \*\*Image available\*\*  
DIGITAL SIGNAL REPRODUCING DEVICE

PUB. NO.: 01-213868 [JP 1213868 A]  
PUBLISHED: August 28, 1989 ( 19890828)  
INVENTOR(s): ARANO YUKARI

ONISHI TAKESHI  
MATSUTANI KIYOSHI  
APPLICANT(s): MITSUBISHI ELECTRIC CORP [000601] (A Japanese Company or Corporation), JP (Japan)  
APPL. NO.: 63-039913 [JP 8839913]  
FILED: February 22, 1988 (19880222)  
INTL CLASS: [4] G11B-020/10  
JAPIO CLASS: 42.5 (ELECTRONICS -- Equipment)  
JAPIO KEYWORD: R101 (APPLIED ELECTRONICS -- Video Tape Recorders, VTR)  
JOURNAL: Section: P, Section No. 964, Vol. 13, No. 524, Pg. 119, November 22, 1989 (19891122)

ABSTRACT

PURPOSE: To ensure the reproducing of PCM data by providing a jump timing pulse generating circuit in a digital signal processing circuit and jumping a corresponding number of addresses to dummy data on the side of the readout of a memory.

CONSTITUTION: The PCM data read out of a track are demodulated by a demodulation circuit 2, and after detecting a symbol value of its header part 4 by a header detecting circuit 3, clocks corresponding to the max. sample number per field are written into a memory 5, whereas data is read out of the memory by a clock relative to a sampling frequency of a sound signal, etc., and outputted to a D/A converter 12 or an interface circuit 13. In this case, addresses of the memory 5 are generated by a DEM side address generating circuit 9, a decoding address generating circuit 8 and a D/A side address generating circuit 7 and selected by a selector 6, and then such addresses on the read side as corresponding to dummy data are jumped over by a pulse generated from the jump timing pulse generating circuit 10, thus omitting the dummy data out of the reproducing PCM data.

18/9/11 (Item 11 from file: 347)  
DIALOG(R)File 347:JAPIO  
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02793026 \*\*Image available\*\*  
WIDE RANGE MOBILE COMMUNICATION SYSTEM

PUB. NO.: 01-090626 [JP 1090626 A]  
PUBLISHED: April 07, 1989 (19890407)  
INVENTOR(s): MORIYAMA KAZU  
APPLICANT(s): KOKUSAI ELECTRIC CO LTD [000112] (A Japanese Company or Corporation), JP (Japan)  
APPL. NO.: 62-245818 [JP 87245818]  
FILED: October 01, 1987 (19871001)  
INTL CLASS: [4] H04B-007/26  
JAPIO CLASS: 44.2 (COMMUNICATION -- Transmission Systems)  
JOURNAL: Section: E, Section No. 791, Vol. 13, No. 323, Pg. 156, July 21, 1989 (19890721)

ABSTRACT

PURPOSE: To attain wide range communication forming a radio communication system even with mixture of digital/analog communication systems by using a cryptographic device to a transmitter-receiver of a radio station so as to realize digital squelch.

CONSTITUTION: The cryptographic device 5 is provided with selection means No.1 and No.2, the means No.1 selects any of radio lines and a most suitable MODEM 53 thereto in each digital station, and the means No.2 switches control modes A-C of a reception signal output. The mode A is the mode operated when all the stations have the device 5 and only when a random number synchronizing signal is detected in the digital reception output of the MODEM 53 in succession to the preamble signal, a decoded

voice is outputted to a speaker 2 as a digital squelch function. The mode B is the mixed communication mode and from which of the digital and analog stations the signal is received is decided depending on the presence of the preamble signal of the reception section and the random number synchronizing signal, and both the stations are made receivable. The mode C is the mode receiving a signal sent from the MODEM 53 without using a random number device 52 and the digital transmission is provided to the line.

18/9/12 (Item 12 from file: 347)  
DIALOG(R)File 347:JAPIO  
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02743947 \*\*Image available\*\*  
COMPRESSED VOICE DATA EXCHANGE SYSTEM

PUB. NO.: 01-041547 [JP 1041547 A]  
PUBLISHED: February 13, 1989 ( 19890213)  
INVENTOR(s): IMAIZUMI YOICHI  
APPLICANT(s): FUJITSU LTD [000522] (A Japanese Company or Corporation), JP  
(Japan)  
APPL. NO.: 62-198384 [JP 87198384]  
FILED: August 07, 1987 (19870807)  
INTL CLASS: [4] H04M-003/00  
JAPIO CLASS: 44.4 (COMMUNICATION -- Telephone)  
JOURNAL: Section: E, Section No. 766, Vol. 13, No. 235, Pg. 142, May  
30, 1989 (19890530)

#### ABSTRACT

PURPOSE: To eliminate need for dividing a digital interface signal line into the talking of its own station and the relay exchange portion by adopting the constitution such that a CODEC and a format conversion section are switched depending whether or not a path is for its own station talking portion or a relay exchange portion so as to eliminate need for an analog interface line.

CONSTITUTION: A path discrimination section 4 is provided to a digital relay exchange 1 and the path deciding section 4 decides whether or not the path is for its own station talking portion or a relay exchange portion and a changeover means 7 of a multiplexer 2 is turned to the position shown in solid lines in case of the own station talking, the signal is through a digital interface signal line 3 and a CODEC 5 to compress the signal at the sender side into a signal rate of 16Kbps or below and to unpack the signal at the receiver side into a rate of 64Kbps for the talking purpose. In case of relay exchange, the changeover means 7 is turned to the position shown in dotted lines, the signal is through a format conversion section 6 at the reception side to bring the signal of 16Kbps or below into a signal of 64Kbps through the addition of a dummy signal, the result is sent to the exchange 1 via the digital interface signal line 3 and exchanged and then sent to the format exchange section 6 of the multiplexer 2 via the digital interface signal line 3, where the dummy signal is excluded and the signal is restored to the original 16Kbps signal and the result is sent to an opposite station.  
? t18/69,k/20,32

18/69,K/20 (Item 6 from file: 350)  
DIALOG(R)File 350:Derwent WPIX  
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0012963821 - Drawing available  
WPI ACC NO: 2003-040953/ 200303  
XRPX ACC No: N2003-032065  
Generating digital bitstream for data rate compression by processing input data blocks to produce shorter blocks with adaptive bit allocation

Patent Assignee: DOLBY LAB LICENSING CORP (DOLB)

Inventor: TRUMAN M M; WATSON M A

Patent Family (3 patents, 98 countries)

Patent Number	Kind	Date	Application Number	Kind	Date	Update	
WO 2002091361	A1	20021114	WO 2002US3705	A	20020208	200303	B
AU 2002235537	A1	20021118	AU 2002235537	A	20020208	200452	E
US 6807528	B1	20041019	US 2001851589	A	20010508	200469	E

Priority Applications (no., kind, date): US 2001851589 A 20010508

#### Patent Details

Number	Kind	Lan	Pg	Dwg	Filing Notes
--------	------	-----	----	-----	--------------

WO 2002091361	A1	EN	15	2	
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National Designated States,Original: AE AG AL AM AT AZ BA BB BG BR BY BZ  
CA CH CN CO CR CU CZ DE DK DM DZ EC EE ES FI GB GD GE GH GM HR HU ID IL  
IN IS JP KE KG KP KR KZ LC LK LR LS LT LU LV MA MD MG MK MN MW MX MZ NO  
NZ OM PH PL PT RO RU SD SE SG SI SK SL TJ TM TN TR TT TZ UA UG UZ VN YU  
ZA ZM ZW

Regional Designated States,Original: AT BE CH CY DE DK EA ES FI FR GB GH  
GM GR IE IT KE LS LU MC MW MZ NL OA PT SD SE SL SZ TR TZ UG ZM ZW

AU 2002235537 A1 EN Based on OPI patent WO 2002091361

#### Alerting Abstract WO A1

NOVELTY - Method consists in capturing blocks of input data and processing them to produce shorter blocks in which some bits allocated from a pool by an adaptive allocation process represent input data, some, which are those remaining in the pool do not and can represent other information, and assembling the shorter blocks to deliver the digital bitstream.

DESCRIPTION - There is an INDEPENDENT CLAIM for a method of processing a digital bitstream.

USE - Method is for low bit rate audio encoding and decoding systems such as \*\*Dolby Digital\*\* and MPEG-2 AAC.

ADVANTAGE - Method uses wasted bits to carry information, so that there is no need to decode and re-encode the bitstream.

DESCRIPTION OF DRAWINGS - The figure shows a \*\*Dolby Digital\*\* encoder.

Title Terms/Index Terms/Additional Words: GENERATE; DIGITAL; BITSTREAM;  
DATA; RATE; COMPRESS; PROCESS; INPUT; BLOCK; PRODUCE; SHORT; ADAPT; BIT;  
ALLOCATE

#### Class Codes

International Classification (Main): G10L-019/00, G10L-021/04

(Additional/Secondary): H04B-001/66

US Classification, Issued: 704229000, 704500000, 370528000

File Segment: EngPI; EPI;

DWPI Class: W04; P86

Manual Codes (EPI/S-X): W04-V10G1J

Alerting Abstract ...USE - Method is for low bit rate audio encoding and decoding systems such as <b>Dolby Digital</b> and MPEG-2 AAC...

...ADVANTAGE - Method uses wasted bits to carry information, so that there is no need to decode and re-encode the bitstream...

#### Original Publication Data by Authority

#### Original Abstracts:

...audio encoding systems, including Dolby Digital and MPEG-2 AAC generate data streams in which unused dummy, fill, stuffing, or null bits exist whenever the bit allocation function in the encoder does not utilize all available bits from a bit pool...

...encoding, a modified encoder may insert information-carrying bits in some or all of the unused bit positions instead of null bits during the encoding process.

...  
...

18/69,K/32 (Item 18 from file: 350)  
DIALOG(R)File 350:Derwent WPIX  
(c) 2007 The Thomson Corporation. All rts. reserv.

0007174277 - Drawing available  
WPI ACC NO: 1995-214940/ 199528  
Related WPI ACC No: 1994-048342  
XRPX ACC No: N1995-168565

EMI suppression coding - psuedorandomises positive and negative going transitions used to represent one binary state in transmitted digital signal

Patent Assignee: TUT SYSTEMS INC (TUTS-N)

Inventor: GRAHAM M H

Patent Family (1 patents, 1 countries)

Patent Number	Kind	Date	Application Number	Kind	Date	Update
US 5422919	A	19950606	US 1992964508	A	19921021	199528 B
			US 1993150451	A	19931110	

Priority Applications (no., kind, date): US 1992964508 A 19921021; US 1993150451 A 19931110

#### Patent Details

Number	Kind	Lan	Pg	Dwg	Filing Notes
US 5422919	A	EN	6	5	Continuation of application US 1992964508

Continuation of patent US 5283807

#### Alerting Abstract US A

The method for communicating binary data uses one binary data state represented by either one of two different signal levels. Another binary data state is represented by another signal level different from the two different signal levels. A pseudorandom signal is generated. One of the two different signal levels is selected as a function of the pseudorandom signal. The binary data is recovered without the use of the pseudorandom signal.

Pref., the pseudorandom signal has a stream of pseudorandom digital signals. A third signal level may be selected to identify the other binary state.

USE/ADVANTAGE - Encoding mechanism using three distinct signal characteristics. Does not require descrambling.

Title Terms/Index Terms/Additional Words: EMI; SUPPRESS; CODE; POSITIVE; NEGATIVE; TRANSITION; REPRESENT; ONE; BINARY; STATE; TRANSMIT; DIGITAL; SIGNAL

#### Class Codes

International Classification (Main): H04L-027/30

US Classification, Issued: 375200000, 380034000, 380046000, 380049000

File Segment: EPI;

DWPI Class: W02

Manual Codes (EPI/S-X): W02-K05A1; W02-K05B1; W02-K05B5

Alerting Abstract ...generated. One of the two different signal levels is

selected as a function of the pseudorandom signal . The binary data is recovered without the use of the pseudorandom signal .

...

...USE/ADVANTAGE - Encoding mechanism using three distinct signal characteristics. Does not require descrambling .

Original Publication Data by Authority

Original Abstracts:

...format. By pseudorandomizing the selection of these transitions, substantial spreading of the spectral energy occurs. Descrambling need not occur since each transition is recognized as the encoded binary state. Thus, the data sender can encode...

Claims:

...for selecting between said two different signal levels under control of said pseudorandom signal; a decoder comprising means for recognizing said two different signal levels as representing said one binary state without the use of said pseudorandom signal ; and , a link coupled between said encoder and decoder.

?

File 2:INSPEC 1898-2007/Sep W5  
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File 6:NTIS 1964-2007/Oct W2  
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File 8:Ei Compendex(R) 1884-2007/Sep W5  
(c) 2007 Elsevier Eng. Info. Inc.  
File 34:SciSearch(R) Cited Ref Sci 1990-2007/Oct W1  
(c) 2007 The Thomson Corp  
File 35:Dissertation Abs Online 1861-2007/Jul  
(c) 2007 ProQuest Info&Learning  
File 65:Inside Conferences 1993-2007/Oct 10  
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File 95:TEME-Technology & Management 1989-2007/Sep W5  
(c) 2007 FIZ TECHNIK  
File 99:Wilson Appl. Sci & Tech Abs 1983-2007/Sep  
(c) 2007 The HW Wilson Co.  
File 144:Pascal 1973-2007/Sep W4  
(c) 2007 INIST/CNRS  
File 256:TecInfoSource 82-2007/Jul  
(c) 2007 Info.Sources Inc  
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File 434:SciSearch(R) Cited Ref Sci 1974-1989/Dec  
(c) 2006 The Thomson Corp  
File 583:Gale Group Globalbase(TM) 1986-2002/Dec 13  
(c) 2002 The Gale Group  
File 56:Computer and Information Systems Abstracts 1966-2007/Sep  
(c) 2007 CSA.  
File 60:ANTE: Abstracts in New Tech & Engineer 1966-2007/Aug  
(c) 2007 CSA.

Set	Items	Description
S1	1434308	FALSE OR INVALID OR NULL OR PSUEDO OR PSEUDO OR RANDOM OR - RAND OR CHAFF OR SEMIRANDOM OR PSUEDORAND? OR PSEUDORAND? OR - DUMMY OR NONCE? ? OR DECOY?
S2	60321	SHILL?? OR SPOOF?? OR PHONEY? ? OR PHONY? ? OR PHONIE? ? - OR FAKE? ? OR SHAM OR SHAMS
S3	37005	S1:S2(1W)(VALUE OR VALUES OR NUMBER? ? OR NUMERAL? ? OR NU- MERIC?? OR ALPHANUMERIC? OR QUANTITY? ? OR QUANTITIES OR CHAR- ACTER? ? OR BIT OR BITS OR DATA OR DATUM? ?)
S4	122794	S1:S2(1W)(DIGIT? ? OR SIGNAL? ? OR PULSE OR PULSES OR INTE- GER? ? OR STRING OR STRINGS OR SUBSTRING? ? OR SEQUENCE OR SE- QUENCES OR SUBSEQUENCE? ? OR NOISE OR VARIABLE? ?)
S5	23306046	USED OR USE OR USING OR USAGE? ? OR EMPLOY? OR UTILIS??? OR UTILISATION? ? OR UTILIZ??? OR UTILIZATION? ?
S6	39140	NONEMPLOY? OR UNEMPLOY? OR NONUTILIS? OR NONUTILIZ? OR UNU- TILIS? OR UNUTILIZ?
S7	20094	NONUSE? OR NONUSING OR NONUSAGE? OR UNUSE? OR UNUSING OR U- NUSAGE? OR DISUSE? ? OR DISUSING OR DISUSAGE?
S8	227500	('NOT' OR WITHOUT OR NEVER OR CANNOT OR T) (1W)S5
S9	4735	(NON OR UN OR DIS)()S5
S10	109	S3:S4(5N)(OMIT? ? OR OMISSION? ? OR EXCLUD? OR EXCLUSION?)
S11	90	S3:S4(5N)S8:S9
S12	161550	DECRYPT? OR UNCRYPT? OR UNENCRYPT? OR DECOD??? OR UNENCOD?- ?? OR UNCOD??? OR UNENCRYPT? OR UNENCIPHER? OR UNENCYIPHER? OR DECIPHER? OR DECPYPER?
S13	52854	UNSCRAMBL? OR DESCRAMBL? OR UNCIPHER? OR UNCYPHER? OR DECO- MPRESS? OR DEPACK? OR UNPACK? OR UNCOMPRESS? OR DECOMPACT? OR UNCOMPACT?
S14	11	S3:S4(5N)S6:S7
S15	1	(S10:S11 OR S14) AND S12:S13
S16	1654	AU=(CHEUNG T? OR CHEUNG, T?)
S17	5	S16 AND S3:S4

S18            5    S17 NOT S15  
S19            2    RD (unique items)

? t15/7

15/7/1        (Item 1 from file: 2)  
DIALOG(R)File 2:INSPEC  
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04777008    INSPEC Abstract Number: C91000799  
Title: Inversion in time (connectionist systems)  
Author(s): Thrun, S.; Linden, A.  
Author Affiliation: Gesellschaft fur Math. und Datenverarbeitung mbH,  
Augustin, West Germany  
Conference Title: Neural Networks. EURASIP Workshop 1990 Proceedings  
p.130-40  
Editor(s): Almeida, L.B.; Wellekens, C.J.  
Publisher: Springer-Verlag, Berlin, West Germany  
Publication Date: 1990 Country of Publication: West Germany    ix+276  
pp.  
ISBN: 3 540 52255 7  
Conference Sponsor: Eur. Assoc. Signal Process  
Conference Date: 15-17 Feb. 1990    Conference Location: Sesimbra,  
Portugal  
Language: English    Document Type: Conference Paper (PA)  
Treatment: Theoretical (T)  
Abstract: Inversion of multilayer synchronous networks is a method which  
tries to answer questions like 'what kind of input will give a desired  
output?' or 'Is it possible to get a desired (output under special  
input/output constraints)?'. The authors describe two methods of inverting  
a connectionist network. Firstly, they extend inversion via backpropagation  
to recurrent, time-delayed and discrete versions of continuous networks.  
The results of inversion is an input vector. The corresponding output  
vector is equal to the target vector except a small remainder. The  
knowledge of those attractors may help to understand the function and the  
generalization qualities of connectionist systems of this kind. Secondly,  
they introduce a new inversion method for proving the nonexistence of an  
input combination under special constraints, e.g. in a subspace of the  
input space. This method works by iterative exclusion of invalid  
activation values. It might be a helpful way to judge the properties of  
a trained network. They conclude with simulation results of three different  
tasks: XOR, morse signal decoding and handwritten digit recognition. (11  
Refs)  
Subfile: c

? t19/7/all

19/7/1        (Item 1 from file: 2)  
DIALOG(R)File 2:INSPEC  
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10477723  
Title: Fractal analysis of yeast cell optical speckle  
Author(s): Flamholz, A.; Schneider, P.S.; Subramaniam, R.; Wong, P.K.;  
Lieberman, D.H.; Cheung, T.D.; Burgos, J.; Leon, K.; Romero, J.  
Author Affiliation: Queensborough Community Coll., CUNY, Bayside, NY, USA  
Journal: Proceedings of the SPIE - The International Society for Optical  
Engineering Conference Title: Proc. SPIE - Int. Soc. Opt. Eng. (USA)  
vol.6084, no.1    p.186-95  
Publisher: SPIE-Int. Soc. Opt. Eng.  
Publication Date: 2006 Country of Publication: USA  
CODEN: PSISDG    ISSN: 0277-786X



SICI: 0277-786X(2006)6084:1L.186:FAYC;1-H  
Material Identity Number: C574-2006-041  
Conference Title: Optical Interactions with Tissue and Cells XVII  
Conference Date: 23 Jan. 2006 Conference Location: San Jose, CA, USA  
DOI: 10.1117/12.646399  
Language: English Document Type: Conference Paper (PA); Journal Paper (JP)

Treatment: Experimental (X)

Abstract: Steady state laser light propagation in diffuse media such as biological cells generally provide bulk parameter information, such as the mean free path and absorption, via the transmission profile. The accompanying optical speckle can be analyzed as a random spatial data series and its fractal dimension can be used to further classify biological media that show similar mean free path and absorption properties, such as those obtained from a single population. A population of yeast cells can be separated into different portions by centrifuge, and microscope analysis can be used to provide the population statistics. Fractal analysis of the speckle suggests that lower fractal dimension is associated with higher cell packing density. The spatial intensity correlation revealed that the higher cell packing gives rise to higher refractive index. A calibration sample system that behaves similar as the yeast samples in fractal dimension, spatial intensity correlation and diffusion was selected. Porous silicate slabs with different refractive index values controlled by water content were used for system calibration. The porous glass as well as the yeast random spatial data series fractal dimension was found to depend on the imaging resolution. The fractal method was also applied to fission yeast single cell fluorescent data as well as aging yeast optical data; and consistency was demonstrated. It is concluded that fractal analysis can be a high sensitivity tool for relative comparison of cell structure but that additional diffusion measurements are necessary for determining the optimal image resolution. Practical application to dental plaque bio-film and cam-pill endoscope images was also demonstrated. (15 Refs)

Subfile: A B

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19/7/2 (Item 2 from file: 2)  
DIALOG(R)File 2:INSPEC  
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08736646 INSPEC Abstract Number: B2003-10-6220W-002

Title: A 0.58-1 Gb/s CMOS data recovery circuit using a synchronous digital phase aligner

Author(s): Cheung, T.S. ; Lee, B.C.

Author Affiliation: Dept. of Network Core Technol., ETRI, Daejeon, South Korea

Conference Title: 2002 45th Midwest Symposium on Circuits and Systems. Conference Proceedings (Cat. No.02CH37378) Part vol.3 p.III-385-8 vol.3

Publisher: IEEE, Piscataway, NJ, USA

Publication Date: 2002 Country of Publication: USA 3 vol.(xlii+686+678+699) pp.

ISBN: 0 7803 7523 8 Material Identity Number: XX-2003-00384

U.S. Copyright Clearance Center Code: 0-7803-7523-8/02/\$17.00

Conference Title: Midwest Symposium on Circuits and Systems

Conference Sponsor: IEEE Circuits & Syst. Soc.; School of Electr. & Comput. Eng. at Oklahoma State Univ

Conference Date: 4-7 Aug. 2002 Conference Location: Tulsa, OK, USA

Language: English Document Type: Conference Paper (PA)

Treatment: Practical (P)

Abstract: A data recovery circuit using a newly proposed synchronous digital phase aligner is realized for multi-link applications. The proposed circuit is implemented with 0.35  $\mu$ m CMOS process technology. The experimental results show that the proposed circuit successfully recovers

incoming 0.58-1 Gb/s of  $2^{31}-1$  pseudo random bit sequence with less than  $10^{-14}$  of bit error rate. (5 Refs)

Subfile: B

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**EAST Search History**

Ref #	Hits	Search Query	DBs	Default Operator	Plurals	Time Stamp
L40	0	((encrypt\$4) and (decrypt\$4) and (original string) and (defin\$4 set of factors) and (encrypt\$4 equation) and (map\$4) and (a set of derivative equations) and (generat\$4 derivative values) and (stor\$4 encrypted string and derivative values) and (provid\$4 false derivatives cannot be used to determine a given factor) and (stor\$4 false derivative values with the generated derivative values) and (us\$4 a set of factor decryption equatins to map) and (decrypt\$4 the encrypted string) and (decryption equation) and (factor mapped through the set of factor decryption equation) and (presence) and (false derivative values with the generated derivative values) and (prevent\$4) and (use with the factor decryption equation to derive the factors in the set of facotrs)).clm.	US-PGPUB; USPAT	ADJ	ON	2007/10/11 22:17

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Best 200 shown

Relevance scale ☐ ☐ ☐ ☐ ☐

# 1 [Cryptography and data security](#)

Dorothy Elizabeth Robling Denning

January 1982 Book

**Publisher:** Addison-Wesley Longman Publishing Co., Inc.

Full text available: [pdf\(19.47 MB\)](#)

Additional Information: [full citation](#), [abstract](#), [references](#), [cited by](#), [index terms](#)

## **From the Preface (See Front Matter for full Preface)**

Electronic computers have evolved from exiguous experimental enterprises in the 1940s to prolific practical data processing systems in the 1980s. As we have come to rely on these systems to process and store data, we have also come to wonder about their ability to protect valuable data.

Data security is the science and study of methods of protecting data in computer and communication systems from unauthorized disclosure ...

# 2 [The relational model for database management: version 2](#)

E. F. Codd

January 1990 Book

**Publisher:** Addison-Wesley Longman Publishing Co., Inc.

Full text available: [pdf\(28.61 MB\)](#)

Additional Information: [full citation](#), [abstract](#), [references](#), [cited by](#), [index terms](#), [review](#)

## **From the Preface (See Front Matter for full Preface)**

An important adjunct to precision is a sound theoretical foundation. The relational model is solidly based on two parts of mathematics: firstorder predicate logic and the theory of relations. This book, however, does not dwell on the theoretical foundations, but rather on all the features of the relational model that I now perceive as important for database users, and therefore for DBMS vendors. My perceptions result from 20 y ...

# 3 [On randomization in sequential and distributed algorithms](#)



Rajiv Gupta, Scott A. Smolka, Shaji Bhaskar

March 1994 **ACM Computing Surveys (CSUR)**, Volume 26 Issue 1

**Publisher:** ACM Press

Full text available:  [pdf\(8.01 MB\)](#) Additional Information: [full citation](#), [abstract](#), [references](#), [citations](#), [index terms](#)

Probabilistic, or randomized, algorithms are fast becoming as commonplace as conventional deterministic algorithms. This survey presents five techniques that have been widely used in the design of randomized algorithms. These techniques are illustrated using 12 randomized algorithms—both sequential and distributed— that span a wide range of applications, including: primality testing (a classical problem in number theory), interactive probabilistic proofs ...

**Keywords:** Byzantine agreement, CSP, analysis of algorithms, computational complexity, dining philosophers problem, distributed algorithms, graph isomorphism, hashing, interactive probabilistic proof systems, leader election, message routing, nearest-neighbors problem, perfect hashing, primality testing, probabilistic techniques, randomized or probabilistic algorithms, randomized quicksort, sequential algorithms, transitive tournaments, universal hashing

#### 4 A graph theoretic framework for preventing the wormhole attack in wireless ad hoc networks

Radha Poovendran, Loukas Lazos  
January 2007 **Wireless Networks**, Volume 13 Issue 1

**Publisher:** Kluwer Academic Publishers

Full text available:  [pdf\(1.37 MB\)](#) Additional Information: [full citation](#), [abstract](#), [references](#), [index terms](#)

Wireless ad hoc networks are envisioned to be randomly deployed in versatile and potentially hostile environments. Hence, providing secure and uninterrupted communication between the un-tethered network nodes becomes a critical problem. In this paper, we investigate the wormhole attack in wireless ad hoc networks, an attack that can disrupt vital network functions such as routing. In the wormhole attack, the adversary establishes a low-latency unidirectional or bi-directional link, such as a wire ...

**Keywords:** geometric random graphs, security, wireless ad hoc networks, wormhole attack


#### 5 Technical reports

 SIGACT News Staff  
January 1980 **ACM SIGACT News**, Volume 12 Issue 1

**Publisher:** ACM Press

Full text available:  [pdf\(5.28 MB\)](#) Additional Information: [full citation](#)

#### 6 A survey of RST invariant image watermarking algorithms

 Dong Zheng, Yan Liu, Jiying Zhao, Abdulmotaleb El Saddik  
July 2007 **ACM Computing Surveys (CSUR)**, Volume 39 Issue 2

**Publisher:** ACM Press

Full text available:  [pdf\(5.53 MB\)](#) Additional Information: [full citation](#), [abstract](#), [references](#), [index terms](#)

In this article, we review the algorithms for rotation, scaling and translation (RST) invariant image watermarking. There are mainly two categories of RST invariant image watermarking algorithms. One is to rectify the RST transformed image before conducting watermark detection. Another is to embed and detect watermark in an RST invariant or semi-invariant domain. In order to help readers understand, we first introduce the fundamental theories and techniques used in the existing RST invariant ...

**Keywords:** Digital image watermarking, Fourier-Mellin transform, ILPM, LPM, RST invariant, Radon transform, feature points, moments, template matching

7 The internet worm program: an analysis

 Eugene H. Spafford  
January 1989 **ACM SIGCOMM Computer Communication Review**, Volume 19 Issue 1

**Publisher:** ACM Press


Full text available:  pdf(2.45 MB) Additional Information: [full citation](#), [abstract](#), [citations](#), [index terms](#)

On the evening of 2 November 1988, someone infected the Internet with a *worm* program. That program exploited flaws in utility programs in systems based on BSD-derived versions of UNIX. The flaws allowed the program to break into those machines and copy itself, thus *infecting* those systems. This program eventually spread to thousands of machines, and disrupted normal activities and Internet connectivity for many days. This report gives a detailed description of the components of the ...

8 An architecture for secure wide-area service discovery

Todd D. Hodes, Steven E. Czerwinski, Ben Y. Zhao, Anthony D. Joseph, Randy H. Katz  
March 2002 **Wireless Networks**, Volume 8 Issue 2/3


**Publisher:** Kluwer Academic Publishers

Full text available:  pdf(365.68 KB) Additional Information: [full citation](#), [abstract](#), [references](#), [citations](#), [index terms](#)

The widespread deployment of inexpensive communications technology, computational resources in the networking infrastructure, and network-enabled end devices poses an interesting problem for end users: how to locate a particular network service or device out of hundreds of thousands of accessible services and devices. This paper presents the architecture and implementation of a secure wide-area Service Discovery Service (SDS). Service providers use the SDS to advertise descriptions of available ...

**Keywords:** location services, name lookup, network protocols, service discovery


9 Security Mechanisms in High-Level Network Protocols

 Victor L. Voydock, Stephen T. Kent  
June 1983 **ACM Computing Surveys (CSUR)**, Volume 15 Issue 2


**Publisher:** ACM Press

Full text available:  pdf(3.23 MB) Additional Information: [full citation](#), [references](#), [citations](#)

10 Magic Functions: In Memoriam: Bernard M. Dwork 1923--1998

 Cynthia Dwork, Moni Naor, Omer Reingold, Larry Stockmeyer  
November 2003 **Journal of the ACM (JACM)**, Volume 50 Issue 6

**Publisher:** ACM Press

Full text available:  pdf(708.05 KB) Additional Information: [full citation](#), [abstract](#), [references](#), [citations](#), [index terms](#)

We prove that three apparently unrelated fundamental problems in distributed computing, cryptography, and complexity theory, are essentially the same problem. These three problems and brief descriptions of them follow. (1) *The selective decommitment problem*. An adversary is given commitments to a collection of messages, and the adversary can ask for some subset of the commitments to be opened. The question is whether seeing the decommitments to these open plaintexts allows the adversary to ...


**Keywords:** Digital signature, Fiat-Shamir methodology, interactive argument, interactive proof system, magic function, selective decommitment, zero knowledge

11 Succinct representation of flexible and privacy-preserving access rights

Marina Blanton, Mikhail Atallah

November 2006 **The VLDB Journal — The International Journal on Very Large Data Bases**, Volume 15 Issue 4

**Publisher:** Springer-Verlag New York, Inc.

Full text available:  [pdf\(525.96 KB\)](#) Additional Information: [full citation](#), [abstract](#)

We explore the problem of portable and flexible privacy preserving access rights that permit access to a large collection of digital goods. *Privacy-preserving* access control means that the service provider can neither learn what access rights a customer has nor link a request to access an item to a particular customer, thus maintaining privacy of both customer activity and customer access rights. *Flexible* access rights allow a customer to choose a subset of items or groups of items ...

**Keywords:** Compact representation, Flexible access rights, Privacy-preserving access rights

12 Paper session DB-10 (databases): query processing 2: Balancing performance and confidentiality in air index

Qingzhao Tan, Wang-Chien Lee, Baihua Zheng, Peng Liu, Dik Lun Lee

October 2005 **Proceedings of the 14th ACM international conference on Information and knowledge management CIKM '05**

**Publisher:** ACM Press

Full text available:  [pdf\(320.99 KB\)](#) Additional Information: [full citation](#), [abstract](#), [references](#), [index terms](#)

Studies on the performance issues (i.e., access latency and energy conservation) of wireless data broadcast have appeared in the literature. However, the important security issues have not been well addressed. This paper investigates the tradeoff between performance and security of signature-based air index schemes in wireless data broadcast. From the performance perspective, keeping low false drop probability helps clients retrieve the information from a broadcast channel efficiently. Meanwhile ...

**Keywords:** indexing techniques, security, wireless data broadcast

13 Symmetric and Asymmetric Encryption

Gustavus J. Simmons

December 1979 **ACM Computing Surveys (CSUR)**, Volume 11 Issue 4

**Publisher:** ACM Press

Full text available:  [pdf\(2.23 MB\)](#) Additional Information: [full citation](#), [references](#), [citations](#), [index terms](#)

14 Credentials: Concealing complex policies with hidden credentials

Robert W. Bradshaw, Jason E. Holt, Kent E. Seamons

October 2004 **Proceedings of the 11th ACM conference on Computer and communications security CCS '04**

**Publisher:** ACM Press

Full text available:  [pdf\(219.13 KB\)](#) Additional Information: [full citation](#), [abstract](#), [references](#), [citations](#), [index terms](#)

Hidden credentials are useful in protecting sensitive resource requests, resources, policies,

and credentials. We propose a significant performance improvement when implementing hidden credentials using Boneh/Franklin Identity Based Encryption. We also propose a substantially improved secret splitting scheme for enforcing complex policies, and show how it improves concealment of policies from nonsatisfying recipients.

**Keywords:** authentication, credentials, identity based encryption, privacy, secret sharing, trust negotiation


## 15 Applying hierarchical and role-based access control to XML documents



Jason Crampton

October 2004 **Proceedings of the 2004 workshop on Secure web service SWS '04**

**Publisher:** ACM Press

Full text available:  pdf(337.19 KB) Additional Information: [full citation](#), [abstract](#), [references](#), [citations](#), [index terms](#)

W3C Recommendations XML Encryption and XML-Digital Signature can be used to protect the confidentiality of and provide assurances about the integrity of XML documents transmitted over an insecure medium. The focus of this paper is how to control access to XML documents, once they have been received. This is particularly important for services where updates are sent to subscribers. We describe how certain access control policies for restricting access to XML documents can be enforced by encryption ...

**Keywords:** XML, encryption, hierarchical access control, role-based access control

## 16 The security of all RSA and discrete log bits



Johan Håstad, Mats Nässtrand

March 2004 **Journal of the ACM (JACM)**, Volume 51 Issue 2

**Publisher:** ACM Press

Full text available:  pdf(311.91 KB) Additional Information: [full citation](#), [abstract](#), [references](#), [index terms](#)

We study the security of individual bits in an RSA encrypted message  $E_N(x)$ . We show that given  $E_N(x)$ , predicting any single bit in  $x$  with only a nonnegligible advantage over the trivial guessing strategy, is (through a polynomial-time reduction) as hard as breaking RSA. Moreover, we prove that blocks of  $O(\log \log N)$  bits of  $x$  are computationally indistinguishable from random bits. The results carry over to ...

**Keywords:** Cryptography, RSA-encryption, bit-security, complexity, discrete logarithms

## 17 New basic technologies for DIM: Pseudonym management using mediated identity-based cryptography



Thibault Candebat, Cameron Ross Dunne, David T. Gray

November 2005 **Proceedings of the 2005 workshop on Digital identity management DIM '05**

**Publisher:** ACM Press

Full text available:  pdf(293.16 KB) Additional Information: [full citation](#), [abstract](#), [references](#), [index terms](#)

Mobile Location-Based Services (LBS) have raised privacy concerns amongst mobile phone users who may need to supply their identity and location information to untrustworthy third parties in order to access these applications. Widespread acceptance of such services may therefore depend on how privacy sensitive information will be handled in order to restore users' confidence in what could become the "killer app" of 3G networks. In this



paper, we present a proxy-based public key infrastructure tha ...

**Keywords:** SEM architecture, identity-based encryption, location-based services, pseudonymity

18 A fuzzy commitment scheme



Ari Juels, Martin Wattenberg

November 1999 **Proceedings of the 6th ACM conference on Computer and communications security CCS '99**

**Publisher:** ACM Press

Full text available: pdf(966.08 KB)

Additional Information: [full citation](#), [abstract](#), [references](#), [citations](#), [index terms](#)

We combine well-known techniques from the areas of error-correcting codes and cryptography to achieve a new type of cryptographic primitive that we refer to as a fuzzy commitment scheme. Like a conventional cryptographic commitment scheme, our fuzzy commitment scheme is both concealing and binding: it is infeasible for an attacker to learn the committed value, and also for the committer to decommit a value in more than one way. In a convent ...

19 Security and correctness: Efficient data protection for distributed shared memory multiprocessors



Brian Rogers, Milos Prvulovic, Yan Solihin

September 2006 **Proceedings of the 15th international conference on Parallel architectures and compilation techniques PACT '06**

**Publisher:** ACM Press

Full text available: pdf(386.29 KB)

Additional Information: [full citation](#), [abstract](#), [references](#), [index terms](#)

Data security in computer systems has recently become an increasing concern, and hardware-based attacks have emerged. As a result, researchers have investigated hardware encryption and authentication mechanisms as a means of addressing this security concern. Unfortunately, no such techniques have been investigated for Distributed Shared Memory (DSM) multiprocessors, and previously proposed techniques for uni-processor and Symmetric Multiprocessor (SMP) systems cannot be directly used for DSMs. T ...

**Keywords:** DSM multiprocessor, data security, memory encryption and authentication

20 Software protection and simulation on oblivious RAMs



Oded Goldreich, Rafail Ostrovsky

May 1996 **Journal of the ACM (JACM)**, Volume 43 Issue 3

**Publisher:** ACM Press

Full text available: pdf(3.44 MB)

Additional Information: [full citation](#), [abstract](#), [references](#), [citations](#), [index terms](#)

Software protection is one of the most important issues concerning computer practice. There exist many heuristics and ad-hoc methods for protection, but the problem as a whole has not received the theoretical treatment it deserves. In this paper, we provide theoretical treatment of software protection. We reduce the problem of software protection to the problem of efficient simulation on oblivious RAM. A machine is oblivious if the sequence in wh ...

**Keywords:** pseudorandom functions, simulation of random access machines, software protection

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